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54 Addressing arrangement for a RAM buffer controller.

57 There is disclosed herein a RAM buffer controller for managing the address input lines of a RAM buffer to simulate the operation of two FIFO's therein. There is also disclosed apparatus for allowing random access by a node processor in a local area network node using the RAM buffer controller to manage transmit and receive FIFO's to have random access to any address in the address space of the buffer without restriction to FIFO boundaries. There is also disclosed apparatus for transmitting packets from said buffer organized into one or two linked lists. Further, there is disclosed apparatus for allowing independent initialization of any of the pointers in the RAM buffer controller which are not currently selected, and for allowing software requests for read or write access by the node processor. Further, there is disclosed apparatus and a method for recording status and length information at the end of a packet instead of in front thereof and for allowing any incoming packet to be flushed without saving status information or to be flushed while saving its status information.

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IMPROVED RBC POINTER ARCHITECTURE

The invention relates to the field of high speed networking applications, and in one aspect, to the field of RAM buffer controllers which control the address ports to random access memory to make the memory appear to be operating as two independent FIFOs.

There are three main application areas for high speed networking systems. The first is in back end networks where the network is used for interconnecting a variety of storage devices such as disks and tape drives with the host computers' backplane. Usually data transfers between disks and CPUs in such configurations call for extremely high throughput. A second application for high speed networks is high performance front end networks wherein a number of computers, terminals, and other devices are connected together by a single transmission medium over which messages are sent between the units on the network. A third application area is in backbone networks where lower speed networks are connected together by high speed networks that act as bridges between the low speed networks. High throughput capability is needed to prevent bottlenecks.

In such networks, a host system, which can be a computer, a terminal, a peripheral, or some other device, utilizes a specially adapted system for interface between the host system and the transmission media. Such systems normally include encoders and decoders for transforming the data into the proper signal format for transmission media and media access controllers for taking care of contentions for the media by different host systems. The host system is connected to a buffer memory by a RAM buffer controller and a data path controller. The purpose of the buffer memory is to store messages which are to be transmitted, until such time as they are transmitted, and to store messages that have been received until such time as the host system is able to retrieve them from memory and process them. The data path controller controls the data inputs of the memory while the RAM buffer controller controls the address inputs of the memory.

In such systems, it is useful to have the buffer memory controlled in such a manner that it simulates the operation of a FIFO memory. This allows the host system to store messages to be transmitted in the buffer memory in a particular order, and allows the message to be transmitted in the same order. The same is true for received messages, i.e., the received message will be stored in buffer memory in a particular order and will be read out of the buffer by the host system in the same order.

There are existing FIFO RAM controllers which are available in the form of integrated circuits. One such circuit is the Signetics 8X60 which manages the address pins of standard off-the-shelf RAMs in such a manner that the RAMs implement a high speed/high capacity FIFO stack. Another such system is the Xilinx Inc. logic cell array in the form of a printer buffer controller model no. XC2064, described in Electronic Design of November 14, 1985 at page 139. Both of these RAM controllers cause standard off-the-shelf RAM to implement a single FIFO buffer.

Unfortunately, neither of these systems has sufficient flexibility and performance in terms of architectural capability to support a complex high speed network application. For example, neither system can cause a single bank of memory to implement two independently operating FIFOs. Further, prior art FIFO RAM controllers generally use counters to keep track of the read and write pointer locations. Such counters cannot generally be initialized to random, desired addresses when any of the other counters are in use in controlling the address applied to the address pins of the controlled memory. The ability to independently set any pointer in the system while another pointer is being used to control the address pins is a useful feature in that it provides great flexibility to the system software designers.

Further, such prior art systems do not generally support random access to any location in memory which access is not constrained by FIFO address limitations. It is useful in implementing various control functions to have a random access capability such that the contents of any address in the memory can be read or changed regardless of whether it is outside currently established FIFO starting or ending addresses.

Another shortcoming of the prior art FIFO controllers is in the area of flushing packets or data which has arrived but which has proven to be not desirable for retention in the memory. Such data may be data which is not addressed to the host system or which may have a transmission error destroying the integrity of the data. The prior art systems do not provide circuitry which allows flushing of data in simple and expeditious fashion.

It is also useful when managing a FIFO to be able to record status and length information associated with each packet in a memory location adjacent to the location as opposed to storing it at the end of the packet. The best location for such status and length information is in the memory location just preceding the first byte of the packet data. Unfortunately the prior art systems do not provide circuitry which adequately supports the ability to simply and quickly write status and length information in the memory location just preceding the first byte of the received packet data.

It is also useful when managing a FIFO to be able to abort a particular packet for some reason without losing its status and length information. The prior art systems do not provide circuitry which easily supports such a control function.

In some networking applications such as token ring networks, it is desirable to be able to transmit data bytes, words, or packets in linked list format, i.e., where the bytes, words, packets to be transmitted are not stored in contiguous locations in memory but instead are spread throughout the memory. Such a linked list has a pointer at the end of each member of the list which points to the starting address for the next member of the list. The prior art systems do not have circuitry which allows linked lists of data to be transmitted. Typically such prior art systems use counters which are incremented by a clock such that the pointers for transmission of data vary continuously from one memory location to the next memory location in contiguous fashion. Such an architecture cannot be easily adapted to linked list transmission.

It is very useful in management of FIFOs to be able to establish flexible starting and ending addresses for the FIFO. The registers which are used to store such starting addresses and ending addresses are not used in linked list transmission. This is because linked list transmission is not a FIFO operation so starting and ending addresses are meaningless. However, during linked list transmission there sometimes arises a need to save the current read address pointer in the middle of transmitting a particular linked list and change the read address pointer to a fixed predetermined address to begin transmitting a particular packet which starts at that address. Such a situation typically arises when, in a token ring situation, a ring recovery packet must be transmitted because of a problem on the network. It is useful to be able to utilize the memory locations which store the starting address and ending address for FIFO operations as new link and save registers for use in vectoring transmission to the starting address of the recovery packet and for saving the current contents of the read pointer.

Finally, it is also useful to be able to send packets in back-to-back fashion with virtually no time between the end of transmission of a first packet and the beginning of transmission of a second packet. Sometimes in FIFO management by a FIFO controller, there are simultaneous access requests by various units in the network interface. For example, the host system may request access to the memory for purposes for storing therein a packet to be transmitted. For such a transaction, the FIFO controller must provide the memory with a write pointer to provide an address into which the first byte of data from the host system will be written. Simultaneously, a packet of data addressed to the host system may be arriving from another system on the network. Such a packet must be stored in the memory as it arrives. In the prior art systems where two FIFO buffers cannot be simultaneously managed, such a situation can never occur, and if it does occur, an error condition will exist. However in systems where such a situation can occur, it will be necessary to arbitrate the simultaneous requests to determine the then current highest priority request. Then, it is useful to have the arbitration occur in a pipelined fashion such that the arbitration is occurring simultaneously with service of the prior request. That is, it is useful to be able to arbitrate a first group of requests and then service the winning request simultaneously with concurrent arbitration with a second group of requests such that upon completion of the service of the winning request from the first arbitration, the service of the winning request from the second arbitration may begin immediately.

Thus a need has arisen for a flexible system for managing external memory so as to provide simultaneous independent FIFO implementation and to satisfy all the above noted needs.

SUMMARY OF THE INVENTION

We will describe a system which can manage the address inputs of a memory such that two independent FIFO buffers can be maintained or, alternatively, a receive FIFO may be implemented while a linked list is implemented for transmission of data. In the preferred embodiment, the invention can support linked list transmission of two or more independent linked lists. In one embodiment, the invention is comprised of a means for storing a plurality of pointer addresses. An output multiplexer is coupled to the first means to select one of the pointer addresses for output to the address pins of the memory. An incrementing circuit coupled to the output of the output multiplexer increments the selected pointer address and supplies it on a BUS which is coupled into the inputs of the first means such that the selected pointer

can be incremented for the next memory access. In the preferred embodiment, the plurality of addresses are stored in a plurality of registers. Each register has an input multiplexer for controlling which of several inputs are selected for coupling into the data input of the register associated with the particular input multiplexer. In one embodiment there is a write address pointer register and a read address pointer register for a transmit FIFO and a write address pointer and a read address pointer for a receive FIFO. In the preferred embodiment, one input of each input multiplexer is coupled to an initialization and another input of each multiplexer is coupled to the output of the incrementing circuit. Thus, the particular pointer which is selected has its input multiplexer switched such that the output of the incrementing circuit is applied to the inputs of the selected pointer and the selected register output is coupled through the output multiplexer to the address pins of the memory. Any other pointer storage register may then be accessed through the initialization and the appropriate input multiplexer such that the contents of that register may be initialized to any desired pointer address.

In another embodiment of the invention, a random access memory pointer storage register is provided. It also has an input multiplexer which has one input coupled to the initialization BUS and another input coupled to the output incrementing circuit.

All of the pointer storage registers and all of the embodiments have one input of their input multiplexers coupled to the output of the associated register. This allows the contents of the pointer register to be maintained at its current value without incrementation as long as the input multiplexer is selecting the input coupled to the register output. This provides maximum flexibility in dealing with external memory since the pointer register contents may be held constant for long periods of time thereby eliminating any timing difficulties which may arise from differing operational speeds of the memory and circuitry which is accessing the memory and allowing pipelined request arbitration as will be described below. The random access memory pointer storage register also shares this capability with the other pointer storage registers.

There is for each of the transmit and receive FIFOs an end address storage register. Both such registers have input multiplexers one input of which is coupled to the initialization bus. This allows the end addresses or the FIFOs to be changed at will. The outputs of the end address registers are switched through an output multiplexer to one input of a comparator, the other input of which is coupled to the output of the output multiplexer coupled to the pointer storage registers. The selected pointer address is constantly compared by the comparator against the appropriate end address for purposes of determining whether the FIFO end boundary has been reached by the currently selected pointer. Control logic coupled to the output multiplexer and input multiplexers insures that the proper end address pointer register is selected for comparison at any particular time based on the currently selected pointer. The output of the comparator goes to the control logic which causes the starting address for the appropriate FIFO to be loaded into the selected pointer when the selected pointer has been incremented to equal the end address. The starting addresses are stored in registers which also have input multiplexers one input of which is coupled to the initialization bus. Thus the starting addresses may be randomly set to any desired address.

The output of the comparator for the end address circuitry is ignored by the control logic when the control logic has selected the random access pointer for output to the address pins of the memory. Thus the random access circuitry of the invention is not subject to any FIFO end address limitations.

Another embodiment of the invention utilizes a shadow write pointer storage register. This register also has an input multiplexer one input of which is coupled to the output of the write pointer storage register for the receive FIFO. Another input of the input multiplexer for the shadow write pointer storage register is coupled to the output of the shadow write pointer register itself. Also one output of the shadow write pointer storage register is coupled to the one input of the input multiplexer for the write pointer storage register for the receive FIFO. The control logic is designed to implement three separate algorithms. They are, respectively: One, an algorithm for reception of normal data packets; two, an algorithm for totally flushing undesired packets; three, an algorithm for aborting received packets without losing the status and length information which is recorded by the algorithm for storing normal received packets.

The algorithm for storing normally received packets allows status and length information to be written in the memory location just preceding the first byte of the received packet. This is true regardless of how long the packet is and whether or not its length is known. This is accomplished by storing the address contents of the write pointer for the receive FIFO in the shadow write pointer register at the beginning of packet reception and skipping the first address location. Packet data is then stored beginning at the newly incremented address stored in the write pointer for the receive FIFO. At the end of packet reception, the address contents of the shadow write pointer register are restored into the write pointer register for the receive FIFO and vice versa and status and length information is written into the location then pointed to by the write pointer. After the status and length information is written, the write pointer register for the receive

FIFO is restored to the address just by following the last byte of data in the received packet by causing the input multiplexers for the write pointer register for the receive FIFO and the shadow write pointer register to select inputs such that the contents of the shadow write pointer register are restored in the write pointer register for the receive FIFO.

5 The control logic to be described is designed also to allow simple and expeditious flushing of undesired packets by causing the contents of the write pointer register for the receive FIFO to be stored in the shadow write pointer register at the beginning of reception of each packet. When a packet is to be flushed, a signal indicates this event to the control logic which then causes the contents of the shadow write pointer register to be stored in the write pointer register for the receive FIFO, thereby restoring the write pointer to the address value marking the beginning of the current packet to be flushed. The control logic for the invention
10 implementation implements the abort algorithm in a similar fashion to the flush algorithm but with some slight variations.

Another embodiment of the invention utilizes a connection of a data bus to an input for the read pointer storage register for the transmit FIFO. When it is desired to transmit linked lists, the control logic causes the
15 input multiplexer for the read pointer for the transmit FIFO to select the input coupled to the data bus. In this manner, the linking address for the next element on the linked list may be loaded into the read pointer for the transmit FIFO such that memory access may be had to the next element on the list through use of this pointer. The control logic causes the output multiplexer to select the read pointer for the transmit FIFO at the output address in such situations.

20 In another embodiment, an additional read pointer for transmission is provided in the form of another read pointer storage register for transmission. This additional register also has an input multiplexer which has one input coupled to the data bus such that two independent linked lists may be transmitted using the two read pointers for transmit operations alternatively or by selecting one pointer for transmission of a first linked list and selecting the second pointer for transmission of the second list after completion of
25 transmission of the first list.

During transmission of linked lists, it sometimes happens that a ring recovery condition will be necessary, and a predetermined packet located at a known location in memory must be transmitted to aid in ring recovery. In such an event, the current contents of the read pointer for the transmit buffer must be saved and a new read pointer for the transmit buffer must be loaded. The new pointer is the address of the
30 fixed packet to be transmitted. The embodiment advantageously uses the storage register for the end address for the transmit buffer for storage of the current read pointer for the transmit buffer. This saving of the current read pointer for the transmit buffer is accomplished through a pathway from the output of the read pointer storage register for the transmit buffer to one input of the input multiplexer for the end address storage register for the transmit buffer. In ring recovery situations, the control logic causes this aforementioned pathway to be selected at the input multiplexer for the end address storage register for the transmit
35 buffer. Simultaneously, the input multiplexer for the read pointer for the transmit buffer is caused to select the output of the starting address pointer storage register to load the new link address. The new link address, i.e., the address of the beginning of the recovery packet, has been previously loaded into the starting address storage register during linked list operations through the initialization bus.

40 In the preferred embodiment, multiple clock cycles are used for each memory access. During one clock cycle of the memory access, the selected pointer is incremented via the incrementing circuit. During clock cycles following incrementation, the selected pointer is caused by the control logic to be maintained at a constant value by causing the input multiplexer to select its input which is coupled to the output of the selected pointer storage register. Thus the pointer can be held at a constant value for as many clock cycles
45 as are needed to complete the transfer. The control logic also simultaneously causes the arbitration logic which is connected to the pointer memory access request lines from the various requesting units to arbitrate the requests during the service cycle or cycles for the winning request from the previous arbitration. In the preferred embodiment, each read or write transfer to the memory is designed as a two clock cycle transfer. During the second clock cycle of the service, the winning request from the first
50 arbitration, the arbitration for the second group of requests, is completed. Thus the winning request for memory access will have been determined by the arbitration circuitry by the completion of the service cycle for the winning request from the previous arbitration. Thus, packets may be transferred in back-to-back fashion without the need for waiting for arbitrations to be completed between packets.

Preferred embodiments of the invention will now be described.

55

BRIEF DESCRIPTION OF THE DRAWINGS.

Figure 1 is a block diagram of a typical system in which the invention might be found.

Figure 2 is a block diagram of a RAM buffer controller of the invention.

5 Figure 3 is a diagram of the physical and logical organization of the two FIFO buffers managed by the RAM buffer controller.

Figure 4 is a more detailed block diagram of one embodiment of the invention.

Figure 5 is a more detailed block diagram of the preferred embodiment of the invention.

Figure 6 is a diagram of the organization in memory of two packets organized as a linked list.

10 Figure 7 is a diagram of the organization of the fields of a typical packet.

Figure 8 is a conceptual diagram of the algorithm used by the invention to receive a normal packet.

Figure 9 is a more detailed diagram of the various multiplexers and registers involved in the three receive algorithms described in the specification.

Figure 10 is a timing diagram of the relationships of the signals involved in the algorithm for receiving
15 normal packets.

Figure 11 is a conceptual diagram of the algorithm for flushing packets.

Figure 12 is a conceptual diagram of the algorithm for aborting packets.

Figure 13 is a timing diagram showing the process for transmitting both linked lists of packets and non-linked packets.

20 Figure 14 is diagram of the memory address locations and their contents used in the example of Figure 13.

Figure 15 is a timing diagram of the pipelined arbitration of service requests carried out by the invention.

Figure 16 is a diagram of an asynchronous system using two systems with DMA needs each having
25 their own clock and which share a common RAM buffer controller which runs on the clock of only one system.

30 DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENT Referring to Figure 1, there is shown a block diagram of a typical system in which the RAM buffer controller of the invention resides. A transmission medium such as a fiber optic cable 20 is used to carry data between systems connected to the medium. A transceiver 22 converts signals on the medium 20 to electrical signals in serial format, and vice versa. Signals to be transmitted on the medium 20 are sent to the transceiver 22 on a bus 24; while signals arriving from the medium are converted to electrical signals of the proper format and output from
35 the transceiver on a bus 26. An encoder/decoder 28 encodes data for transmission arriving on a bus 30 into the proper code selected for transmission on the medium 20. Data arriving from the medium 20 is decoded in the encoder/decoder 28 and output on a bus 32.

40 Every local area network has a protocol. Further, every network has a topology, such as star, ring, or other topologies which may be developed in the future. To provide the flexibility in architecture to accommodate all such differences in topology and protocol, the protocol or topology-specific functions are implemented in the media access controller circuit (FORMAC) 34. The design of the media access controller, encoder/decoder, and transceiver are known for the various topologies and protocols which today exist. Further, they are not critical to the invention, since the RAM buffer controller of the invention is
45 designed for use in any topology and any protocol with any type of medium. Thus, for a particular type of topology and protocol, a particular design for the media access controller will be selected. The same is true for the transceiver and for the encoder/decoder.

All the circuitry of Figure 1 is designed to service the communications needs of a host system 36. This host system may be any one of a number of different types of logic systems, such as a computer, disk
50 drive, personal computer, work station, or other peripheral. This host system may be big or small, and it may be fast or slow.

The host system 36 has a need to transmit data packets to the other host systems on the network and to receive data packets therefrom. Because the host system may be busy doing other tasks when a receive packet arrives on the medium 20 addressed to the host system 36, it is necessary to temporarily store the
55 received packet until the host system 36 has the opportunity to process it. Likewise, the host system 36 has a need from time to time to transmit packets to other host systems on the network. However, the network can accommodate only one pair of transmitting and receiving host systems at any particular time. Therefore, when the host system 36 wishes to transmit a packet, it is sometimes found that its access to

the medium 20 is blocked by virtue of another pair of host systems' having exclusive control of the medium 20. Therefore, it is sometimes necessary to temporarily store the packet to be transmitted in a buffer memory until such time as the medium 20 is granted to the host system 36. The temporary storage functions mentioned herein are performed by a buffer memory 38. This buffer memory has an address port
 5 connected to an address bus 40 for receiving addresses which control the particular location in the buffer memory 38 from which data is retrieved or to which data is written. Data retrieved or data to be written is transmitted on or received over a data bus 46.

A data path controller 43 is coupled to the buffer memory 38 via the data bus 46 and serves to control the data transmission to and from the buffer memory 38 on the data bus. The RAM buffer controller of the
 10 invention is symbolized by the block 44, which is coupled to the buffer memory 38 by the address bus 40. The purpose of the RAM buffer controller is to arbitrate requests for access to the buffer memory 38 by the various elements in the system of Figure 1, select one request to be honored, and to output an address pointer to the buffer memory 38 on the address bus 40 to cause the buffer memory 38 to output data from or store data in the proper address. The RAM buffer controller of the invention maintains a set of pointer
 15 addresses stored in registers therein which can cause the buffer memory 38 to simulate the operation of two independently operating buffer memories. In the preferred embodiment, these two simulated buffers are FIFO organized, but in other embodiments, such as embodiments outside the local area network field where the RAM buffer controller is being used for stack operations, these two buffers may be LIFO organized such that the read and write pointers are managed or otherwise controlled such that a last in, first
 20 out operation is performed.

The host system 36 can send data to the buffer memory 38 or receive data from the buffer memory 38 on a data bus 46. This data bus 46 is also connected to a data input of the RAM buffer controller. Further, the data bus 46 is connected to a data port of a node processor.

The node processor 52 is a typically a microprogrammed or conventional microprocessor-based system
 25 used for overseeing the operation of the RAM buffer controller 44, the data path controller 43, and the media access controller 34. The node processor is coupled to these three units by an NPI bus 54. The main function of the node processor is to initialize the contents of various registers in the RAM buffer controller 44 and to initialize the data path controller and the media access controller 43 and 34, respectively. The node processor also responds to various system level and packet level interrupts. In
 30 simple systems the node processor may be the minimum state machine required for the system requirements. In complicated systems, the node processor may be a sophisticated enough machine to execute all the higher level protocols in the 7 layer ISO model. The node processor does its work with the RAM buffer controller and the other units of the system using the NPI bus 54 and various bus handshake and instruction lines. The other lines in Figure 1 which have not been discussed so far are control and
 35 handshake signals which are interchanged between the units to control the direction of data transfer on the various buses, acknowledge the receipt of certain signals, signal the occurrence of various events, and otherwise control the operation of the system. The particular control signals which are relevant to the invention will be discussed later herein, as the need arises. The control lines 56 and 58 between the node
 40 processor and host system implement a handshake which is user dependent and depends upon partitioning functions between the host and the node processor. For instance, in large systems the host may have a powerful and complex node processor 52 which is used for off loading various networking-specific chores.

The node processor can talk to the buffer memory 38 by making a node processor service request on the control bus 101 to the RAM buffer controller and transmitting data to or receiving data from the buffer
 45 memory 38 on the data bus 46. A typical node processor would usually consist of a microprocessor with assorted peripheral chips for performing DMA, interrupts, etc. It would also have some local memory, and it would treat the RAM buffer controller and the data path controller as peripherals for networking functions. The node processor receives status information from the RAM buffer controller, the data path controller, the media access controller, and the buffer memory 38. The node processor therefore has complete control
 50 over and knowledge of the state of these units.

The node processor runs either synchronously or asynchronously with respect to the network clock.
 55 The network clock is derived from the incoming data by the transceiver 22. That is, the transceiver interfaces to the medium and derives clocking information from the encoded data stream and passes this clocking signal on to the media access controller. The media access controller then typically converts the serial data stream to 8-bit parallel data for output on the bus 32. The data path controller 43 then converts the 8-bit parallel data from the media access controller and encoder/decoder to 32-bit parallel data for
 transmission on the data bus 46.

Referring to Figure 2, there is shown a more detailed block diagram for the RAM buffer controller 44. The RAM buffer controller (hereafter RBC) includes request arbitration logic 62, which serves to arbitrate requests from the node processor, the data path controller (sometimes hereafter referred to as the DPC) and the host system for access by way of a DMA transfer to or from the buffer memory 38. DPC requests for access to the memory are made using handshake signals directly exchanged between the RBC and DPC on the bus 64 in Figure 1. These access requests by the DPC are routed by the handshake logic on line 78 to the request arbitration logic which arbitrates the requests and acknowledges them on line 79. The acknowledge signals sent back to the DPC on bus 64 in response to the acknowledge signals on line 79 will be discussed later. The handshake with the DPC is performed by DPC handshake logic 66 in Figure 2.

Each of the node processor, host system, and DPC can request either a read access or a write access to the buffer memory 38. The signal NPRDQ to the arbitration logic 62 represents a node processor read request, while the signal NPWRRQ represents a node processor write request. There is an acknowledge signal on the status signal bus 68 for the reception of each of the aforementioned signals. The signal HSRDRQ represents a host system read request, and the signal HSWRRQ represents a host system write request. Again each of the host system service requests has an acknowledge signal on the status signal bus 68. The DPC handshake logic 66 receives the signal DRDREQS and DRDREQA. These signals represent, respectively, a read request synchronous and a read request asynchronous from the DPC. The signals DRDACKS and DRDACKA, respectively, represent the acknowledge signals for the aforementioned DPC read request signals. The signal DWRREQ represents a data path controller write request, and the signal DWRACK represents the acknowledge signal for that write request.

There is an RBC control logic collection of circuitry which is represented by blocks 70, 72, 66 and 96 in Figure 2. This control logic in later figures will be collectively referred to as control logic 125. The RBC control logic component 70 and the RBC control logic component 72 serve to control the various multiplexers in the system (not shown in Figure 2) based on the particular mode in which the RBC is operating at any particular time. The mode that the machine is operating in is determined by the state of the control signals entering the system from all sources. The RBC control logic 70 also generate control signals which are transmitted to the buffer memory 38 on the bus 41. These control signals consist of the signal CSO NOT, WR NOT, and RD NOT. The signal CSO NOT is a chip select signal which is sent to the buffer memory 38 to select that chip, thereby causing its address and data ports to become active. This signal is low at any time that a read or write memory operation is taking place. The signal WR NOT tell the buffer memory that the data at its data input port coupled to the data bus 42 should be stored in the address then existing on the address bus 40. The signal RD NOT tells the buffer memory 38 to retrieve the data stored at the address then existing on the address bus 40 and output it onto the data bus 42.

The addresses to the buffer memory 38 are stored in a plurality of register symbolized by the memory address pointer block 80 in Figure 2. Depending on which access request is currently selected for service, one of these pointers is selected for output on the address bus 40. This current address is then incremented by the incrementer block 82 and fed back to the input of the pointer register selected via the bus 84. Thus, as long as that pointer register is selected, its contents will be continuously incremented according to a service request pattern established by the RBC control logic in synchronization with a byte clock signal BCLOCK on a line 86.

In certain modes of operation, specifically, the linked list mode, it is necessary to input address data into certain of the registers of the RBC as pointers to the next item on the list. This address information is input on the data bus DBI 42. In this mode of operation, the RBC control logic causes a multiplexer 86 to select the DBI bus 42 such that its address data can be loaded into the proper registers. The multiplexer 86 symbolically represents a collection of input multiplexers which will be defined more specifically in later drawings herein.

A boundary comparator 88 compares the value of the currently selected pointer on bus 90 against the end address pointers stored in the memory boundary pointer registers 92, and generates a control signal on a line 94 coupled to the RBC control logic 70 to indicate when the currently selected pointer has reached the end address of the particular FIFO of interest.

An instruction decoder 96 decodes instructions received on an instruction bus 98 from the node processor and generates control signals on the bus 100 coupled to the RBC control logic 70 and the incrementer 82 to cause the proper events to occur to carry out an instruction. Table 1 below is a listing of the instructions in the RAM buffer controller instruction set and indicates what each instruction does. An NPI initialization bus 102 from the node processor carries initial address values established by the node processor for the starting and ending addresses of the receive and transmit FIFO buffers in the buffer memory 38. It also carries initial address values which can be loaded into most of the pointer registers in

the RAM buffer controller. Bus interface logic 104 receives a chip select signal CSI NOT (active low as are all signals where the signal name is followed by NOT) from the node processor indicating the RAM buffer controller is selected for an operation. Another interface signal DS NOT is active low when data on the NPI bus 102 is synchronous to the byte clock signal BCLOCK on the line 86. In such a case, the DS NOT signal behaves much like a data strobe on slave devices connected to a standard microprocessor. There is a signal INSTR4 in the instruction bus 98 which serves to signal whether a read or write transaction of the registers in the RAM buffer controller is requested. Each register in the RAM buffer controller can be read through a multiplexer not shown in Figure 2. When the signal INSTR4 and CSI NOT are low, a write operation is in progress from the node processor to the RAM buffer controller on the NPI bus 102, and the signal DS NOT should become active whenever the data to be written is valid. The DS NOT signal should stay active and the data on the NPI bus should remain valid until the RAM buffer controller forces the signal READY NOT to go low. In case of a read operation, when INSTR4 is high and CSI NOT is low, the node processor makes the signal DS NOT go low whenever it is ready to accept the read data. When the signal CSI NOT is high, the signals on the bus 98 are ignored, and the NPI bus 102 is kept in the three-state mode.

Table I
Instruction Set

	INST	Decimal value	Inst. mnemonic	Inst. function
	43210			
10	00000	0	IRESET	Software Reset
	00001	1	ILDRPXS	Load Synchronous RPX
	00010	2	ILDRPXA	Load Asynchronous RPX
	00011	3	ILDRPR	Load RPR
	00100	4	ILDMAR	Load MAR
15	00101	5	ILDWPX	Load WPX
	00110	6	ILDWPR	Load WPR
	00111	7	ILDSAX	Load SAX
	01000	8	ILDSAR	Load SAR
	01001	9	ILDEAX	Load EAX
20	01010	10	ILDEAR	Load EAR
	01011	11	ILDCMD	Load Command
	01100	12	INPRDNI	Node processor read request, no increment
	01101	13	INPWRNI	Node processor write request, no increment
25	01110	14	INPRDWI	Node processor read request with increment
	01111	15	ICLDSTS	Clear Dynamic Status
	10000	16	INPWRWI	Node processor write request with increment
30	10001	17	IRDRPXS	Read Synchronous RPX
	10010	18	IRDRPXA	Read Asynchronous RPX
	10011	19	IRDRPR	Read RPR
	10100	20	IRDMAR	Read MAR
35	10101	21	IRDWPX	Read WPX
	10110	22	IRDWPR	Read WPR
	10111	23	IRDSAX	Read SAX
	11000	24	IRDSAR	Read SAR
	11001	25	IRDEAX	Read EAX
40	11010	26	IRDEAR	Read EAR
	11011	27	IRDCMD	Read Command
	11100	28	IRSV4	Reserved
	11101	29	IRSV5	Reserved
	11110	30	IRDSTAT	Read Status
45	11111	31	IRDDSTS	Read Dynamic Status

A command and status interrupt logic 106 generates the nonmaskable interrupt signal NMINTR NOT. This signal is used for interrupting the node processor when the fatal error sets the status bit in the RBC.

Referring to Figure 3, there are shown both a physical and logical organization diagram of the receive and transmit FIFOs contained within the buffer memory 38. On the left is the physical memory organization. Two pointers, SAR and EAR, define the starting and ending addresses, respectively, of the receive FIFO. Another nonoverlapping block of memory is used for the transmit FIFO, which has its starting address defined by the pointer SAX and its ending address defined by the pointer EAX. On the right in Figure 3 is the functional organization of these two receive and transmit FIFOs. For the receive FIFO, there is a read pointer, RPR, which defines the current location for read operations by any requesting unit. There is also a write pointer, WPR, which defines the current address location for any write operations in the receive FIFO. For the transmit FIFO, the write pointer, WPX, defines the current location for any write operations. Likewise, the read points, RPX, defines the current location for any read operations from the transmit FIFO.

For the receive FIFO, write operations correspond to data arriving on the medium 20, which is then written into the buffer memory 38. Read operations for the receive FIFO correspond to accesses by the node processor or the host system to retrieve the data from the buffer memory that has arrived on the medium 20. For the transmit FIFO, write operations correspond to the host system, or possibly the node processor, writing data into the buffer memory 38 for later transmission on the medium 20. Read operations on the transmit FIFO correspond to the process of taking data to be transmitted out of the buffer memory 38 and converting it to the proper code to be transmitted on the medium 20.

Referring to Figure 4, there is shown a more detailed block diagram of one embodiment of the RAM buffer controller of the invention. For purposes of clarity of explanation, a general description of the operation of the circuitry of Figure 4 will be given, without reference to timing details at this point, to enable the reader to grasp the general operation of the system. More details will be given later on the operation of the system in connection with the description of timing diagrams detailing various modes of operation.

To transmit a packet, the host system must write the packet to be transmitted into the buffer memory 38 in Figure 1. To do this, the buffer memory 38 must be supplied with an address or addresses in the transmit FIFO in which the data is to be stored. To initiate this process, the host system will assert the signal HSWRRQ in Figure 2 to the request arbitration logic 62. There may be other service requests pending at the time that the aforementioned service request by the host system is made. It is the job of the arbitration logic to select the highest priority currently pending request for service. The priority of service requests is that DPC request have the highest priority, followed in order by node processor requests and host system requests.

When the host system write request is selected for service by the arbitration logic 62, the arbitration logic transmits a select signal to a G multiplexer having inputs coupled to the outputs of the pointer storage registers in the RAM buffer controller indicating that the write pointer for the transmit buffer, WPX, is to be selected for output. This causes the output of the WPX pointer storage register to be coupled to the address bus 40 by the G multiplexer. Prior to this time, the node processor will have issued a number of instructions on the instruction bus 98 to cause loading of the starting and ending addresses for at least the transmit FIFO, and the initial addresses for at least the write pointer storage register for the transmit FIFO. These instructions are decoded by the instruction decode logic 96 and cause the proper ones of the multiplexers A-F and H-I to select their inputs coupled to the NPI bus 102 at the time that the pointer storage register to be loaded has valid data waiting for it on the NPI bus 102. For example, the node processor will load the starting address for the transmit FIFO into the SAX pointer storage register during one stage of the initialization process. Accordingly, the node processor will issue the instruction ILDSAX (see Table 1 herein), indicating that the node processor wishes to load the SAX pointer storage register. When this instruction is decoded, the A multiplexer will select its number 2 input (hereafter all inputs for multiplexers in the drawings will be referred to by their input number when counting from left to right) such that the NPI bus is coupled to the data input 106 for the SAX register. The proper starting address can then be loaded on the NPI bus and will be clocked into the SAX register for storage on the next BCLOCK upward transition. The other pointers detailed in Figure 3 are loaded by issuing one of the instructions shown in Table 1 to cause the particular multiplexer coupled to the data input of each pointer's storage register to select the NPI bus for coupling to the data input of the pointer storage register. The desired address data can then be loaded into the pointer storage register using the NPI bus.

When the G multiplexer receives a select signal indicating the host write request has been honored for service, this multiplexer will select the WPX register output coupled to its second input for coupling to the address output bus 40. The control logic 70 will also receive the select signal via the line 76 indicating which request is currently being serviced. The control logic will then generate control signals on the bus 41 to control the buffer memory. Specifically, the control logic 70 will assert the signals CSO NOT and WR NOT to tell the buffer memory that a write operation is requested. Further, the control logic will generate control signals on a bus 43, which will cause the D multiplexer to select its first input for coupling to the data input 108 of the WPX pointer storage register. The first input of the D multiplexer is coupled to the output of the incrementer circuit 82. The input of the incrementer circuit is coupled to the address bus 40 such that the incrementer always receives as its input the currently selected pointer address. The incrementer, which is a conventionally designed ALU, increments the WPX pointer by 1 and couples it to the first input of the D multiplexer. In some embodiments, each address for a selected pointer will be valid for one period of the byte clock signal and will be incremented during the next period. In such embodiments, the control logic then causes the input multiplexers C-F for the selected pointer to select its first input coupled to the output of the incrementer such that every clock cycle will cause the loading of an incremented address into the pointer storage register currently selected. This goes on for as many clock cycles as the selected pointer is selected. In the preferred embodiment, as will be described more fully

below, each pointer address is valid for two clock cycles of the byte clock signal. This allows pipelined arbitration as will be described in more detail below. Thus, in the preferred embodiment, when the WPX pointer is selected for output, the control logic causes the D multiplexer to select its number 4 input during the two clock cycles of the bus transfer period on the data bus for one word, and causes the D multiplexer to select its number 1 input at the end of the second clock cycle of the service period so that the WPX pointer will be incremented by 1 for the next two cycle service period if the WPX pointer is still selected during the next service period. This process continues until all the words of the packet to be stored in the buffer memory by the host system 36 are so stored.

When the host system is writing data to the buffer memory 38, the select signal from the arbitration logic 62 to the G multiplexer causes that multiplexer to select the WPX pointer and output it on the address bus 40. A comparator 88 has one of its comparison inputs coupled via the line 90 to the address bus 40. The other input of this comparator 88 is coupled to the output of the J multiplexer, which has as its inputs the end address pointers for both the transmit and receive FIFO buffers. These addresses have been previously loaded by the node processor 52 in the EAX and EAR pointer storage registers. This loading was accomplished as described above by causing the H multiplexer to select its number 2 input coupled to the NPI bus after issuance of a load EAX instruction, and by causing the I multiplexer to select its number 1 input after issuance of a load EAR instruction by the node processor. When the host system is writing data to the buffer memory, the control logic, which will be collectively referred to by reference number 125 in Figure 4, causes the J multiplexer to select its number 1 input for coupling to the B input of the comparator 88. Thus, the comparator 88 compares the selected pointer value, in this case WPX, to the end address for the transmit buffer. The comparator 88 outputs a signal on a line 94 coupled to the control logic 125 which indicates whether the selected pointer is equal or not equal to the appropriate end address for the buffer to which the selected pointer pertains, i.e., the transmit buffer in this example. The control logic allows incrementing of the selected pointer to proceed until such time as the signal on line 94 indicate that the selected pointer is equal to the end address. At that point, the control logic 125 causes the input multiplexer for the selected pointer, in this case the D multiplexer, to select the input coupled to the output of the appropriate starting address pointer storage register, i.e., the number 2 input for the D multiplexer. This causes the storing of the starting address for the appropriate buffer, in this case the transmit buffer, into the selected pointer storage register, i.e., the WPX pointer storage register. Thus, incrementing for further transactions begins from the starting address of the appropriate buffer. To summarize the process for the example of the host system writing data into the buffer memory 38, the WPX pointer will be continuously incremented as long as it is selected until such time as it equals the end address for the transmit buffer. At that time, the signal on the line 94 will change states, and the control logic 125 will cause the D multiplexer to select its number 2 input, thereby loading the starting address for the transmit buffer from the SAX pointer storage register into the WPX pointer storage register. If the WPX pointer continues to be selected for further transactions, incrementing will continue from the starting address of the transmit buffer.

The data path controller and media access controller modules will be signaled that there is a packet ready for transmission on the medium 20 said packet having been loaded into the buffer memory 38 by the host system or by the node processor. When access to the medium 20 is obtained, the data path controller 43 will assert a service request requesting that data be read from the buffer memory 38. This request from the data path controller takes the form of assertion of the signal DRDREQS or DRDREQA in Figure 2. These two signals stand for data path controller read request synchronous and data path controller read request asynchronous. The assertion of a data path controller read request causes the control logic 125 to generate the proper select signal to the G multiplexer to select the RPX pointer. This pointer is the read pointer for the transmit buffer. The control logic 125 also causes the C multiplexer to select the number 3 input during the two cycles of the service request so as to maintain the RPX pointer constant at its then existing value. At the end of the second cycle of the service request, the control logic 125 causes the C multiplexer to select its number 1 input coupled to the output of the incrementer 82 to allow the RPX pointer to be incremented. As long as the RPX pointer is selected, it will be maintained constant during two clock cycles and will be incremented at the end of every service cycle. That is, the address or RPX pointer output on address bus 40 causes the buffer memory 38 to retrieve the data at the specified address and output it on the data bus 42. The data path controller and media access controller circuitry then processes this data and, with the help of the encoder/decoder 28 and transceiver 22, transmits the data on the medium 20.

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As in the case of the WPX selection, the RPX pointer is constantly compared to the end address for the transmit buffer by the comparator 88 while RPX is selected. If at any time RPX become equal to EAX, the signal on line 94 changes state, thereby causing the control logic 125 to generate a select signal to the C multiplexer. This select signal causes the C multiplexer to select its number 4 input, thereby loading the starting address for the transmit buffer into the RPX pointer storage register from the SAX pointer storage register.

When data arrives on the medium 20, the data path controller will assert a write request in the form of activation of the signal DWRREQ. This causes the control logic 125 to generate a select signal to the G multiplexer to select its number 4 input, thereby outputting the WPR pointer on the address bus 40. As before, the control logic 125 will cause the F multiplexer to select its number 3 input for the two clock cycles of the transaction time for one word of the incoming data. At the end of the two cycle service period, the control logic 125 will cause the F multiplexer to select its number 1 input, thereby allowing the WPR pointer to be incremented at the start of the two clock cycles representing the next WPR address and service period. Again, the WPR pointer address on the address bus 40 is compared by the comparator 88 to an end address for the appropriate buffer, i.e., the receive buffer. In this case the control logic will order the J multiplexer to select its number 2 input such that the end address for the receive buffer is coupled to the B input of the comparator 88. When the WPR pointer equals the EAR end address, a signal on the line 94 will change states, thereby causing the control logic 125 to generate a select signal to the F multiplexer, causing it to select its number 4 input, thereby loading the starting address for the receiver buffer into the WPR pointer storage register.

After the incoming data has been loaded into the buffer memory 38, the host system 36 may, at its leisure, assert a read request signal to the arbitration logic 62. This event would take the form of assertion of the signal HSRDRQ on bus 103 in Figure 2. If this is the highest priority request then pending, the arbitration logic 62 would generate a select signal which would cause the G multiplexer to select its number 3 input, thereby outputting the RPR pointer address on the address bus 40. The control logic 125 would be notified on the control line 76 of the winning request and would know that a read transaction was requested. The control logic 125 would then generate the proper control signals on the bus 41 to the buffer memory to select the buffer memory and cause it to perform a read transaction. The buffer memory would then access the data stored in the address pointed to by the address on bus 40 and outputted on the data bus 42. The host system could then read the data on the data bus 46 and do with it whatever needed to be done. The control logic 125 would also generate the proper control signals to cause the E multiplexer to select its number 3 input during the two cycles of the read transaction for each word and to select its number 1 input at the end of each service period to allow the RPR pointer storage register contents to be incremented by one.

The comparison of the RPR pointer to the end address for the receive buffer EAR occurs as described above for the other pointers. If the RPR pointer becomes equal to EAR, the control logic 125 orders the E multiplexer to select its number 4 input, thereby loading the starting address for the receive buffer into the RPR pointer. As in the case of the other pointers, the E multiplexer selects its number 3 input for the two cycles of the service period, and it selects its number 1 input at the end of the service period to allow the RPR pointer to be incremented.

That completes the basic description of the operation of the RAM buffer controller in managing two FIFO buffers in a single physical memory 38.

INDEPENDENT INITIALIZATION

It is important for flexibility of design and use of a RAM buffer controller to be able to initialize any of the pointers of the system independently of any of the other pointers of the system. To that end, the invention utilizes a series of input multiplexers denoted by the letters C-F in Figure 4, each of which has one input which is coupled to the NPI bus 102. The output of the incrementer 82 is coupled to another input of each of these input multiplexers. Because the NPI bus and the output of the incrementer on the bus 83 do not have to pass collectively through a common multiplexer with its output coupled to an input of each of the input multiplexers, it is possible to initialize any of the pointer storage registers from the NPI bus while selecting another pointer for output on the address bus 40. That is, if any particular pointer address storage register is selected for output of its contents by the G multiplexer onto the address bus 40, then any of the other pointer storage registers not so selected may be initialized through the NPI bus 102 without interfering with the selection and incrementation of the pointer address in the selected pointer storage register. This is

done by issuing the appropriate load instruction on the instruction bus 98 thereby causing the control logic 125 to cause the appropriate input multiplexer to select its input coupled to the NPI bus. If the node processor attempts to initialize the pointer currently selected, an error condition results and an interrupt is generated.

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RANDOM ACCESS POINTER FOR NODE PROCESSOR SERVICE REQUESTS

In Figure 4 there is a pointer storage register designated MAR which is used to store a pointer address for servicing node processor read or write request. The random access pointer MAR can be initialized by the NPI bus in the same manner as any of the other pointers. To initialize the MAR, the node processor issues a load MAR instruction which is decoded and causes the control logic 125 to send a select signal to the K multiplexer causing it to select its number 2 input to couple the NPI bus to the data input of the MAR pointer storage register. The desired address can then be loaded on the NPI bus, and it will be stored in the MAR register. Thereafter the node processor may assert a service request in the form of activating one of the signals NPRDQ or NPWRRQ in Figure 2. If either of these lines is activated, the request arbitration logic 62 causes the G multiplexer to select its number 5 input, thereby outputting the MAR pointer on the address bus 40. Depending on which of the node processor signals is activated, the control logic 125 generates the proper control signals on the bus 41 to cause the buffer memory 38 to either read data stored at the address on address bus 40 and output it on the data bus 46, or input data from the data bus 46 and store it in the address existing on the address bus 40.

When the MAR random access pointer is selected for output by the G multiplexer, the comparator 88 compares it to the ending address for the transmit buffer or for the receive buffer, but the output signal of the comparator on line 94 is ignored by the control logic 125. The reason for this is that the random access pointer is not limited by any address limitations of the FIFO or LIFO buffers managed by the RAM buffer controller. That is, the random access pointer may be set to any address within the physical address space of the buffer memory 38.

Any node processor service request has the second highest priority in arbitration, and is serviced only if a data path controller request is not present simultaneously. There is a command register in the control logic 125 which contains a command bit called CENNPRQ. This command bit can be set or cleared to enable or disable the node processor requests. It would normally be set to disable node processor requests while a DPC request was being serviced. When a node processor request is being serviced, the control logic generates an acknowledge signal which is active for the duration of the transfer. There is a separate acknowledge signal for a node processor read request and another acknowledge signal for a node processor write request. The node processor inputs data from the buffer memory on the data bus 46 at its data port in Figure 1 and the same data port is used for outputting data to be stored in the buffer memory 38. Since the random access pointer register is not subject to reloading with a starting address upon reading an ending address of either the transmit or receive FIFOs, the random access pointer MAR will be continuously incremented during the entire time it is selected until it reaches its highest value of all ones. At that point it will be reset to a value of all zeros.

Node processor service requests are typically used by the node processor to do block DMA transfers of data to or from the memory. The acknowledge signals for the node processor service request are output from the control logic 125 on a bus 68 in Figure 4. A status signal NPPND is provided on the bus 68 to indicate to the node processor whether there are any requests pending to be serviced. Likewise, when a host system service request is being serviced, one of two acknowledge signals on the bus 68 dedicated to acknowledging host service requests is active. There is a separate acknowledge signal for a host read transaction and another acknowledge signal for a host write transaction. Both types of host requests have equal priority and are serviced alternatively. The same is true for node processor requests in that node processor read and write requests cannot be serviced simultaneously, but must be serviced alternately. Two command bits are provided to independently enable or disable the host system read request and host system write request input ports. These command bits are resident in a command register (not shown) located in the control logic 125.

Typically the random access pointer MAR is used by the node processor to look at headers of packets after a packet has been received. That is, the random access pointer can be used as part of the mechanism used by the node processor to filter out packets that are not addressed to the host system 36 or which have been received but in which CRC errors have been detected.

The node processor has a way of preventing the incrementer 82 from incrementing the MAR pointer such that this pointer can be used for a read-modify-write cycle on buffer memory 38, which otherwise would not support this type of data transaction. To implement a read-modify-write transaction or otherwise block the incrementer 82 from incrementing the selected pointer, the node processor would send an instruction via the instruction bus 98 in Figure 2 to the instruction decoder 96. This instruction could be one of the reserved instructions in Table 1. The instruction could be encoded and a signal transmitted on the line 100/101 to the incrementer 82 to prevent it from incrementing. In the case where the G multiplexer has selected the MAR pointer, the node processor could use the initial address of the MAR pointer and assert a read request. The data output from the buffer memory could then be input by the node processor 52 and modified. The node processor could then assert a write request and transmit a do not increment instruction to the RAM buffer controller. Then the node processor would issue a write request and then write the modified data to the same memory address from which it was read. The do not increment signal on line 101 in Figure 4 is used to block the carry-in input to the incrementer ALU to block incrementation.

In the preferred embodiment, there are four instructions given in Table 1 for allowing the node processor to assert read or write requests either with or without incrementing. These instructions are given in Table 1 above with the mnemonics INPRDNI, INPWRNI, INPRDWI and INPWRWI. Each of these instructions can be asserted on the instruction bus 98 by the node processor, and each is decoded by the instruction decoder 96. When each instruction is decoded, one of the four line in the bus 105 coupled to the output of the instruction decoder is asserted. These read and write request lines are OR'ed in the request arbitration logic 62 with the hardware read and write request lines on the bus 101 from the node processor. Thus the node processor can assert its DMA requests either by asserting one of the signals NPRDQ or NPWRRQ on the bus 101 or by issuing one of the software instructions noted above. However, the node processor cannot assert both a hardware request and a software request at the same time since the MAR is selected by the control logic for any node processor service request and inconsistencies could develop.

The foregoing instructions may be used to implement a read-modify-write sequence as follows. The node processor issues instruction INPRDNI to read the memory location pointed to by the MAR but the MAR is not incremented. The data is then modified, and the node processor issues instruction INPWRWI to write the modified data back into the same memory location which is then pointed to by the MAR. The MAR is then incremented during the a portion of this write access.

The foregoing instructions may also be used to implement a write-verify sequence as follows. The node processor issues the instruction INPWRNI to write data into the location pointed to by the MAR, and the MAR is not incremented. Thereafter, the node processor issues the instruction INPRDWI to read the contents of the location just written and the MAR contents are incremented during this read transaction.

DATA CHAIN FOR LINKED LIST PACKET TRANSMISSION

It is useful in local area network design to be able to transmit packets which are not located contiguously in memory. That is, it is desirable to have the flexibility to be able to put packets for transmission into random locations in memory and to send them in a predetermined order, regardless of their physical locations in memory. To do this requires a linked-list organization of the data packets which are stored in memory and a mechanism to follow the pointer addresses.

Referring to Figure 6, there is shown such a linked list organization for a linked list of packets which can be transmitted with the use of the RAM buffer controller of the invention. Each packet has a byte called a descriptor at the beginning of the packet followed by one or more bytes of data. The descriptor byte has transmit status and length information stored therein. Each packet is followed by a pointer byte which contains the starting address of the next member of the list to be transmitted. The pointer byte 130 is the pointer byte for packet number 1, while the pointer byte 132 is the pointer for packet number 2. The descriptor 134 is the descriptor for packet number 2. If packet number 2 were to be transmitted following the transmission of packet number 1, the pointer 130 would point to the address of the descriptor 134.

The circuitry of the DPC contains a circuit to identify when a word read from the buffer memory 38 is a pointer address. When the pointer 130 is read at the end of transmission of packet number 1, the contents of pointer 130 would be loaded in the RAM buffer controller as the new value for the RPX read pointer for the transmit buffer. On the next read request for access to the transmit buffer, the address of the descriptor 134 would be found in the RPX pointer storage register and transmission of packet number 2 would begin from that point. If, on the other hand, packet number 2 was to be transmitted first and packet number 1 was

to be transmitted following the transmission of packet number 2, then the contents of pointer 132 would be the address of the descriptor 128 for packet number 1. Thus, when descriptor 132 was read, its address contents would be loaded into the RPX pointer storage register, and upon a read request for the transmit buffer again occurring, transmission of packet number 1 would start at the location of the descriptor 128.

5 Referring to Figure 5, there is shown a block diagram of the RAM buffer controller in its preferred embodiment. The portion of the RAM buffer control circuitry which supports the linked-list transmission of packets is the data bus 42 and its connection to one of the inputs of the input multiplexer controlling the data inputs of the read pointer for the transmit buffer. In the preferred embodiment, there are two read pointer storage register for the transmit buffer, designated, respectively, RPXS and RPXA. The data bus 42
10 is coupled to one input of each input multiplexer controlling the data inputs for the RPXS and RPXA pointer storage registers.

When the data path controller is performing read operations from the transmit buffer during linked-list transmit mode, it detects the reading of pointer addresses such as pointers 130 and 132. When such a pointer is read by the data path controller 43 in Figure 1, the data path controller asserts an interface signal
15 to the RAM buffer controller called LDRPXS or LDRPXA, depending upon which pointer storage register is assigned to the particular linked list being transmitted. When the control logic detects the activation of this signal on bus 64, it signals the appropriate one of the input multiplexers C or L to select the input coupled to the DBI bus 42. Because at that time the contents of the pointer will be resident on the DBI bus 42, the pointer address is loaded into the appropriate one of the RPXA or RPXS pointer storage registers on the
20 next upward transition of the BCLOCK signal.

The embodiment of Figure 5 shows two full/empty comparators 140 and 142. The comparator 140 has one compare input coupled to the output of the RPXS pointer and the other input coupled to the output of the WPX pointer. The comparator 140 serves to compare these two pointers and generate a transmit buffer full/empty signal called RPXEQWPX, which is coupled to the control logic 125. The signal RPXEQWPX
25 become active when the read pointer and the write pointer for the transmit buffer become equal. This condition indicates either that the transmit buffer is full or that it is empty. The control logic 125 determines which case is true by examining the last transaction which occurred in the service request prior to the cycle during which the signal RPXEQWPX became active. If the last transaction before the two pointers became equal was a read, the transmit FIFO is empty. If the last transaction before the two pointers became equal
30 was a write, then the transmit FIFO is full.

The comparator 142 acts in a similar manner as the comparator 140 in monitoring the full/empty status of the receive FIFO. The comparator has one compare input coupled to the output of the RPR register, and another compare input coupled to the output of the WPR register. When the read pointer for the receive FIFO and the write pointer for the receive FIFO become equal, the comparator 142 activates the signal
35 RPREQWPR coupled to the control logic 125. The rule for determining whether the receive FIFO is full or empty is the same as the rule for the transmit FIFO.

When the linked-list mode of transmission of packets, the output signals from the comparators 140, 142 and 182 are ignored by the control logic 125. The reason for this is that in the linked-list mode of transmission, the RPXS and RPXA pointers will be jumping around the physical address space of the
40 memory. Thus, the comparison of these pointer values to an end address value and full/empty comparisons will be meaningless.

UTILIZATION OF SAX AND EAX AS NEW LINK AND SAVE POINTER STORAGE REGISTERS

45 While in the linked-list mode of transmission, there is not starting address or ending address for the transmit FIFO. Accordingly, the SAX and EAX pointer storage registers are being wasted. There are some situations such as ring recovery, where it is necessary to jump to a fixed packet location and transmit that packet over the medium to aid in recovery of the ring from some error conditions. To effectively do this, it
50 is necessary to load the starting address of the ring recovery packet into the RPXS pointer storage register as the "new link" pointer. Thus, when the ring recovery signal or other new link signal is asserted, the control logic 125 causes the C multiplexer in Figure 5 to select its number 1 input in Figure 5 so as to load the current contents of the SAX pointer storage register into the RPXS storage register. Before this occurs, however, it is necessary for the node processor to load the new link address in the SAX pointer storage
55 register. This is done as described earlier for initialization of the pointers in the RAM buffer controller.

After the ring recovery has occurred, it is useful to resume transmission of the packet which was interrupted by the ring recovery situation at the location where transmission ceased. To do this, the contents of the RPXS pointer storage register must be saved as it exists at the time that the new link pointer is loaded into the RPXS pointer storage register. The EAX pointer storage register can be used for this purpose. To do this, however, a new pathway must be added between the output of the RPXS storage register and the input of the EAX pointer storage register. This pathway is bus 150 in Figure 5. This bus 150 is coupled to one input of the H multiplexer. At the time that a ring recovery situation or other "new link" situation arises, the control logic sends a select signal to the H multiplexer causing it to select its number 2 input coupled to this bus 150 such that the contents of the RPXS pointer storage register is loaded into the EAX pointer storage register on the next clock uptick. After this loading occurs, the control logic 125 sends a select signal to the H multiplexer causing it to select its input number 3, which is coupled to the output of the EAX storage register, such that the contents of the EAX pointer storage register is maintained constant.

The node processor can read the contents of any pointer storage register in the RAM buffer controller of Figure 5 by sending a read instruction for the particular pointer storage register of interest to the control logic 125 over the instruction bus 98. There is a multiplexer which has one input coupled to the data output of each of the pointer storage registers in the RAM buffer controller. In Figure 5 this multiplexer is shown at 152.

When it is desired to resume packet transmission in the linked-list mode after a ring recovery situation has occurred, the node processor may send a read EAX instruction to the control logic 125 which will be decoded and will cause the read multiplexer 152 to select its input coupled to the output of the EAX pointer storage register. The data in the EAX pointer storage register may then be reloaded into the RPXS pointer storage register via the NPI bus 102. The data path controller can then again assert a read request for access to the transmit buffer such that RPXS is selected by the G multiplexer for output on the address bus 40. Transmission will then resume with the packet which was interrupted by the ring recovery situation.

In the preferred embodiment, in a ring recovery situation, the node processor does not restore the contents of the EAX pointer storage register into the RPXS pointer storage register. Instead, it uses the contents of the EAX pointer storage register as an indication of how many packets were successfully transmitted before the interruption. The node processor then figures out which packet must be the first one sent from the linked-list transmission that was interrupted, i.e., it figures out and loads into RPX the address of the descriptor of the first packet which was not successfully transmitted.

RECEIVING NORMAL PACKETS AND FLUSHING AND ABORTING RECEIVED PACKETS

Figure 7 shows a typical data structure for a packet transmitted on the medium 20 in Figure 1. A first field or group of fields will hold control information, and these fields will be followed by a destination address indicating the particular host system to which the data portion of the packet is addressed. The source address follows the destination address and gives the network address of the originating host system which transmitted the data. Following the source address there will be one or more bytes of data comprising the subject matter of the packet followed by one or more bytes of CRC data which provide check bits to improve the reliability of the data.

When a packet is transmitted on the network, all node processors on the network begin copying the incoming data into the buffer memory. Typically, there may be ten or more control bytes in the control field before the destination address arrives. These ten or more control bytes must be copied into the buffer memory 38 in case the destination address proves to be the address of the particular host system at that node. If, however, after the control bytes have been stored, the destination address arrives and is compared by the node processor to the address of the host system at that node and the two do not match, then it is necessary to flush all the control bytes and other bytes received in that packet so as not to waste memory space with unneeded data. Also, it sometimes happens that entire packets are received but, upon processing the CRC bytes which follow the data portion of the packet, it is discovered that an error in transmission has occurred and the data in the packet is not reliable. In such a case, it is again necessary to flush all the bytes of the packet to free memory space for valid packets.

The mechanism in the preferred embodiment to do this flushing is the shadow write pointer storage register, designated SWPR in Figure 5. This shadow write pointer has an M multiplexer which controls the data inputs of the SWPR pointer storage register. This M multiplexer has one of its inputs coupled to the output of the WPR storage register and has the other input coupled to the output of the SWPR pointer storage register. The output of the SWPR pointer storage register is also coupled to one of the inputs of the F multiplexer, which controls the input of the WPR pointer storage register.

The concept of packet flushing involves saving the contents of the WPR into the SWPR pointer storage register at the beginning of reception of each packet. In this way the SWPR storage register stores the contents of the WPR pointer storage register as it existed at the time when packet reception commenced. Then, if it is necessary to flush the packet, the contents of the SWPR register are restored to the WPR pointer storage register, thereby resetting it to the position that it had at the beginning of the reception of the faulty packet.

To best understand the process for flushing packets it will be instructive to first examine the process of receiving a normal packet. Figure 8 is a conceptual diagram of the algorithm for receipt of a normal packet. Figure 9 is a more detailed block diagram showing the various multiplexer select signals for the F and M multiplexers controlling the WPR pointer storage register and SWPR pointer storage register. These select signals are manipulated by the control logic to implement the algorithm of Figure 8 in accordance with the timing diagram shown in Figure 10. Figure 10 relates the timing relationship of all of these select signals shown in Figure 9 to implement the algorithm of Figure 8 for reception of a normal data packet. Referring jointly to Figures 8, 9, and 10, the first step in receiving a normal data packet is to initialize the WPR pointer storage register to a valid address. To do this, the control logic 125 asserts the signal LDNPWPR at 170 in Figure 10. This causes the F multiplexer to select its B input for application to the data inputs of the WPR pointer storage register. The B input is connected to the NPI bus, which will have previously been driven by the node processor to an address state of an address in which the node processor wishes the first word of the incoming packet to be stored. If packets are to be stored contiguously in memory, the initialization step may be omitted after the first initialization occurred, since the WPR pointer storage register will be continuously incremented until it reaches the end address of the receive FIFO. After a sufficient amount of time for the loading of the WPR register to occur, the signal LDNPWPR is deactivated at 172. The WPR pointer storage register must then be held at the initialization value. Therefore, the control logic asserts the signal HLDWPR at 174 in Figure 10. The signal HLDWPR remains asserted for as many clock cycles as occur prior to the beginning of reception of an incoming packet. In other words, reception of a packet does not actually begin until step 2 in Figures 8 and 10.

The beginning of receipt of a packet is signaled by the assertion of a signal BRCVPKT on bus 64 in Figure 2 from the data path controller. Upon the assertion of the signal BRCVPKT, several things happen. Conceptually, as seen from step 2 of Figure 8, the contents of the WPR pointer storage register are loaded into the shadow write storage register, SWPR, and the WPR pointer storage register is incremented by one, thereby skipping the first location after saving the contents of the WPR. To implement this sequence of events, the control logic deactivates the select signal HLDWPR, thereby deselecting the D input of the F multiplexer. As long as the D input of the F multiplexer is selected, the WPR output data will be fed back into its D input such that every uptick of the clock signal BCLOCK will reload the WPR with its presently existing data, thereby holding it at a constant address.

To load the contents of the WPR pointer storage register into the SWPR pointer storage register, the control logic at t1 asserts the signal LDWPSWPR at 178 in Figure 10. At 180 the control logic deactivates the signal HLDWPR, which is normally active at all times except times when the SWPR pointer storage register is being loaded with new data. The select signal LDWPSWPR causes the M multiplexer in Figure 9 to select its input A. This selects the output of the WPR pointer storage register for application to the D input of the SWPR pointer storage register. Upon the next occurrence of a BCLOCK upward transition, the transfer will occur. To cause the WPR pointer storage register contents to be incremented by one, the control logic at time t1 in Figure 10 asserts the select signal LDINCWPR, as shown at 182 in Figure 10. The control logic at time t1 will assert the signal SELWPR. This causes the G multiplexer to select the input coupled to the WPR pointer storage register output for output on the address bus 40. This couples the WPR pointer storage register current contents to the input of the incrementer 82 in Figure 5. The incrementer 82 then adds one to the contents of the WPR pointer storage register and outputs the incremented value on the bus 83. The assertion of the signal LDINCWPR at t1 in Figure 10, as shown at 182, causes this incremented value present at the C input of the F multiplexer in Figure 9 to be loaded into

the WPR pointer storage register on the next upward transition of the BCLOCK signal. The result is that the WPR pointer storage register contents are incremented by one after storing the original contents in the shadow pointer register. After all these events have transpired, at a time t2 the various signals activated and deactivated at time t1 revert to the state they had just prior to the occurrence of time t1.

5 Between times t2 and t3 in Figure 10 some arbitrary number of BCLOCK periods will have passed before the first bytes of data in the packet being received first appear on the data bus 46 in Figure 1.

At time t3 the first data packet word arrives. The occurrence of this event is signaled by the assertion of the signal DWRREQ in Figure 2 from the data path controller 43. This represents a data path controller write request. It causes the control logic 125 to continuously increment the WPR pointer storage register after each word is written and to continue to hold the shadow write pointer storage register at the value loaded from the WPR pointer storage register in step 2 of Figure 8. During the time between t2 and t3, the WPR pointer storage register was placed in the hold status by assertion of the signal HLDWPR, as shown at 186 in Figure 10. This caused the F multiplexer to select its D input for application to the data input of the WPR pointer storage register. At time t3 this hold is removed, and the control logic asserts the signal LDINCWPR, as shown at 188, and asserts the signal SELWPR, as shown at 190. At the same time, the control logic asserts the signal CSO NOT, as shown at 192, to enable the buffer memory 38, and the control logic asserts the signal WR NOT at 194 to inform the buffer memory that a write transaction is requested. The assertion of SELWPR causes the current content of the write pointer storage register, WPR, to be output on the address bus 40. The first word of the received packet to be written into the buffer memory at the address on bus 40 has previously or simultaneously been placed on the data bus 46 by the data path controller 43. The buffer memory then writes the data on the data bus into the address pointed to by the current contents of the WPR pointer storage register. The simultaneous assertion of LDINCWPR causes the contents of the WPR pointer storage register to be incremented by one on the next upward transition of the BCLOCK signal by virtue of the incrementer having the WPR pointer applied to its inputs and the incremented value applied to the C input of the F multiplexer.

At time t4 all the signals asserted or deactivated at time t3 revert to their states just prior to time t3.

Between times t4 and t5 any number of BCLOCK periods occur, and the control logic asserts the signal HLDWPR to hold the contents of the WPR pointer storage register constant. At time t5, the data path controller again asserts the write request signal DWRREQ and places another word from the packet onto the data bus 46 in Figure 1. The same sequence of events that occurred at t3 now occurs again at t5, and the word then on the data bus 46 is stored in the memory location in the buffer memory just following the memory location in which the word stored between times t3 and t4 was stored.

When the data path controller and the other circuitry in the system of Figure 1 detect the end of the incoming packet, the data path controller asserts the signal ERCVPKT on bus 64 in Figure 2. This is the signal which causes the exchange of pointer register contents symbolized by step 4 in Figure 8 to occur. Conceptually, the contents of the shadow write pointer storage register are loaded into the WPR pointer storage register and the contents of the WPR pointer storage register are loaded into the shadow write pointer storage register simultaneously. The control logic 125 implements this swap by asserting and deasserting various signals at time t7 in Figure 10. As shown at 196, the select signal LDSWWPR is asserted, causing the F multiplexer to select its E input for application to the data input of the SWPR pointer storage register. The E input of the F multiplexer is coupled to the output of the SWPR pointer storage register. Therefore, this causes loading of the then existing contents of the SWPR register into the WPR pointer storage register. To do this successfully, the control logic deasserts the select signal HLDWPR as shown at 198 to release the hold on the contents of the WPR pointer storage register. The control logic also asserts the select signal LDWPSWPR, as shown at 200, to cause the M multiplexer to select its A input for application to the data input of the SWPR pointer storage register. The A input of the M multiplexer is coupled to the output of the WPR pointer storage register, and therefore, the contents of the WPR pointer storage register are loaded into the SWPR register upon the next upper transition of the BCLOCK signal. The control logic also deactivates the select signal HLDSWPR at that time to allow the swap to occur.

By the time t8 in Figure 10, the contents of the WPR pointer storage register will have been swapped with the contents of the SWPR pointer storage register. This means that the contents of the WPR pointer storage register are now equal to the address that the WPR register stored at the beginning of reception of the packet in step 2 of Figure 8 before the WPR register was incremented in step 2. That is, if a packet occupies address locations 1-100 in memory, then the data words in the packet will occupy locations 2-100 and the WPR pointer storage register will contain an address pointing to storage location 1 upon completion of step 4 in Figure 8. In other words, location 1 is the memory location skipped by the incrementation that occurred in step 2 of Figure 8.

Further, location 1 is the location in which the status and length information for the packet is stored. It is useful in local area network design to store status and length information with each packet. This information indicates whether or not it is a valid packet and the length of the packet in terms of the number of memory locations consumed by the packet. It is useful to put this status and length information at the beginning of the packet rather than at the end of the packet. The reason for this is that it simplifies the software overhead in the node processor, because the node processor knows exactly where to look for the status and length information in each packet without having to know how long the packet is, which it would have to know if the status and length information were stored at the end of the packet.

Step 5 in Figure 8 represents the step of writing the status and length information in the first memory location of the packet. In the hypothetical example just given, this would mean the status and length information would be written into memory location 1. This occurs at time t_9 in Figure 10.

In Figure 10 the interval between t_9 and t_{10} shows the control signals as they are activated and deactivated to implement both steps 5 and 6 of Figure 8. Step 5 is implemented by asserting the signal SELWPR at 204, which outputs the current contents of the WPR to the address bus 40. Also asserted are the signals CSO NOT and WR NOT. These two signals activate the buffer memory and tell it that a write transaction is about to occur. The buffer memory then writes into the pointed to location the data on the data bus, which, at time t_9 , is the status and length information.

Step 6 in Figure 8 is implemented in Figure 10 by the assertion of the signal LDSWWPR at 206 and the deactivation of the signal HLDWPR at 208. The events at 206 and 208 in Figure 10 cause the F multiplexer to select the D input, thereby loading the contents of the SWPR back into the WPR pointer storage register. This leaves the WPR pointer storage register with a pointer address pointing to the memory location which is one greater than the last memory location in which a word from the just received data packet is stored. In the hypothetical example given herein, the WPR pointer storage register would now be pointing to address location 101.

Next, in step 7 in Figure 8, a zero bit is written into a specific bit location of the memory location just following the last word in the data packet just received, and the WPR pointer storage register is held constant at that pointer address value. That is, in the hypothetical example, step 7 represents the step of writing a zero into a particular bit of memory location 101 and holding the WPR pointer at 101. The reason that this is done is to cause a zero bit to be written into a specific bit of the memory location just following the data packet. When the node processor or host system is reading data out of the receive buffer and it reads locations 101, it will see this particular bit set at zero. This will tell the node processor or host system that all that follows is not a valid packet. This simplifies the software record-keeping functions of the node processor or the host system, eliminating the need to keep track of the number of valid packets and their lengths, which are currently stored in the receive buffer. Valid packets are designated by writing a one into the designated bit during step 5 of Figure 8 when the status and length information is written into the first memory location of the packet. Step 7 is represented at t_{10} in Figure 10 by the assertion of the signal caps HLDWPR at 210 to hold the WPR content steady by forcing the F multiplexer to select its D input and by continued selection of the WPR output by the G multiplexer by continued assertion of the signal SELWPR at 212. Writing of the zero bit, as required by step 7 of Figure 8, is accomplished by the continued assertion of the signals CSO NOT and WR NOT at 214 and 216 between times t_{10} and t_{11} .

PROCEDURE FOR FLUSHING RECEIVED PACKETS

Referring to Figure 11, there is shown a conceptual diagram of the algorithm for flushing packets which, for one reason or another, need not be retained in buffer memory. Steps 1 through 3 of this algorithm are identical to steps 1 through 3 of the algorithm for receiving a normal packet shown in Figure 8. Upon detection of the fact that a packet is not addressed to the host system at the particular node or that a CRC error has occurred, the DPC asserts a signal FSHRCVP, meaning flush this packet. This signal, FSHRCVP, is one of the interface signals to the RBC control logic 72 in Figure 2, which comes from the DPC over the bus 64. When this signal is asserted, the contents of the SWPR pointer storage register are loaded into the WPR pointer storage register. This restores the WPR to the pointer address which it had at the time when packet reception started. In the hypothetical example given above, step 4 in Figure 11 would restore the WPR pointer address to memory address 1. Since no status and length recording step will have been performed in the algorithm at the time that FSHRCVP is asserted, the status bit at memory location 1 will still be a zero from the previous packet reception. That is, if the previous packet received was normal, the system in step 7 of the algorithm to receive normal packets will have written a zero bit into the memory

location following the last word in the previous packet. In the hypothetical example, this assumes that the previous packet's last word will be stored in memory location 0. This zero status bit will indicate to the node processor or host system that all the data which is in a memory location following memory location 1 is invalid data thereby effectively flushing the packet.

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PROCEDURE FOR ABORTING RECEIVED PACKETS

The algorithm of Figure 12 is the process which the RAM buffer controller implements in the event that a packet is to be aborted. Steps 1 through 3 of the aborted packet algorithm are the same as the first three steps of the algorithm for receiving normal packets as previously described, and will not be described again here. When a packet is aborted, the WPR pointer storage register is reset to the memory location at which it pointed at the time that packet reception began, and status and length information is written. This transaction is symbolized by step 4 in Figure 12 wherein the contents of the SWPR pointer storage register are reloaded into the WPR pointer storage register. Because in step 2 the contents of the WPR pointer storage register were saved into SWPR, this restores WPR to the pointer address it had at the time when packet reception began. In the hypothetical example given above, the WPR pointer storage register would then be pointing to address location 1.

Step 5 of Figure 12 represents the step of writing the status and length information for the aborted packet. To do this, the control logic would assert the select signal SELWRP, forcing the G multiplexer to select the output of the WPR register. Further, the control logic would assert the chip select signal and the write signal to the buffer memory to enable the buffer memory and tell it to write the data then existing on the data bus 46 into the memory pointed to by the current contents of the WPR, i.e., address location 1. The node processor, or DPC, will have loaded the status and length information on the data bus prior to or simultaneously with the chip select and write signals to the buffer memory. Step 5 also represents the process of incrementing the WPR pointer storage register by one. This is done in the same manner as described above with reference to Figure 10. In the hypothetical example, the WPR register would be pointing to address location 2 at the completion of step 5 of Figure 12. Step 6 of Figure 12 is the final step in the aborted packet algorithm. This step involves writing a zero status bit into the predetermined bit in the memory location following the status and length information for the aborted packet. In the hypothetical example, this zero status bit would be written into memory address location 2. The manner in which this is done is identical to step 7 in the algorithm represented by Figure 8.

35 PROTOCOL FOR LINKED AND NONLINKED PACKET TRANSMISSION

Referring to Figure 13, there is shown a timing diagram for the sequence of events which occurs during both linked-list transmission and nonlinked-list transmission.

Nonlinked-list transmission occurs as follows. To load the buffer memory with the packet to be transmitted, the host system asserts a host system write request in the form of assertion of the signal HSWRRQ, shown at 220. This assertion is sensed by the priority arbitration logic 62 in Figure 2 and acknowledged by assertion of the signal HSWRACK at 222. This acknowledgement signal is sent back to the host system via the status bus 68 in Figure 2. A hypothetical example will be used to illustrate the operation of the system, as shown in Figure 13. In this hypothetical, the contents of the WPX pointer storage register at the outset of the transmission is zero, as are the contents of the RPX pointer storage register.

When the host system write request is acknowledged at 222, the G multiplexer is ordered by the control logic to select the output of the WPX pointer storage register. This causes the zero stored in the WPX pointer storage register to be applied to the input of the incrementer 82. The incrementer then outputs a one on line 83. In the preferred embodiment, the host system write request, acknowledge signal lasts for two clock cycles, i.e., between the times t2 and t4 in Figure 13. In other embodiments, one clock cycle might be used, or more than two clock cycles might be used. The incrementer output line 83 stays constant at one, however, from the time of the transition 222 to the transition 224 at t4.

The transition 224 causes the control logic 125 in Figure 4 to cause the D multiplexer in Figure 4 to select its number 1 input, thereby loading the one at the incrementer output into the WPX pointer storage register. During the time from t2 to t4 in Figure 13, the first word of the packet to be transmitted was loaded into the buffer memory at address location zero.

At time t4 the host system asserts another write request, as shown at 226. This causes an acknowledgement signal by the host system by assertion of the signal HSWRACK, as shown at 228. The transition at 228 causes the arbitration logic 62 to issue a select signal to the G multiplexer which selects the output of the WPX pointer storage register for connection to the address bus 40. This applies the current contents of WPX to the incrementer input, which then increments the one currently residing in the WPX register to a two, and outputs it on line 83. During the time from t5 to t7, the second word to be transmitted is loaded by the host system into the buffer memory 38 at address one via the data bus 46.

At time t7, the host system asserts another write request, as shown at 230. Simultaneously, the data path controller asserts a read request by activating the signal DRDREQS on bus 64 in Figure 2. This assertion is shown at 232 in Figure 13. At this point the control logic 125 in Figure 4 blocks requests from the node processor or host and awards priority to the data path controller, because the data path controller always has the highest priority. Basically, the process that is occurring here is the host system is trying to fill up the transmit buffer with words from a packet that it wishes to transmit. Substantially simultaneously, the data path controller is starting to empty words from the buffer memory and transmit them out on the medium. The award of priority to the data path controller is shown by the assertion of the data path controller read request acknowledge signal, DRDACKS, as shown at 234 in Figure 13. The award of priority to the data path controller precludes sending an acknowledgement signal to the host system to acknowledge its write request shown at 230. This is shown by the absence of assertion of the acknowledge signal HSWRACK between times t8 and t9, as shown at 236. The assertion of the data path controller acknowledge signal, DRDACKS, at 234 causes the control logic to send a select signal to the G multiplexer, causing it to select the contents of the RPX pointer storage register for output on the address bus 40. This supplies the current contents of the RPX pointer storage register on address bus 40 at time t8 to the buffer memory 38. Thus a zero pointer address would be supplied to the buffer memory for the DPC read operation commencing at time t8 in the hypothetical example used here. Accordingly, the buffer memory would output the contents of the zero address location between times t8 and t10. Figure 14 shows a hypothetical set of memory addresses and their contents. Accordingly, a 100 would be placed on the data bus in response to the DPC read request corresponding to the word written into the memory by the host system between times t2 and t4. Also the RPX value of zero would be applied to the input of the incrementer, which would then increment it and output a one on line 83 at time t8.

The deactivation of the DPC acknowledge signal at time t10 at 236 causes the G multiplexer in Figure 4 to select its input coupled to the output of the incrementer for coupling to the data input of the RPX pointer storage register. This loads the one at the output of the incrementer into the RPX pointer storage register, thereby incrementing this pointer from zero to one.

Upon completion of servicing of the DPC read request at time t10, the arbitration logic 62 acknowledges the pending host system write request, asserted at time t7. This is done as shown at 238 in Figure 13 by asserting the host system write acknowledge signal HSWRACK at time t10. The assertion at 238 causes the G multiplexer to once again select the contents of the WPX pointer storage register for output on the address bus. The incrementer then takes the current contents of the WPX pointer storage register at address two and increments it to address number three. When the host system write request service is completed at 240 and time t12, the signal HSWRACK is deactivated, which causes the control logic to order the D multiplexer to select its number 1 input coupled to the output of the incrementer, thereby loading the three at the output of the incrementer as the new incremented pointer address for the WPX pointer storage register.

At time t11, the data path controller asserts another read request, as shown at 242. This DPC read request is acknowledged at time t12 by the arbitration logic 62, as shown by the assertion of the signal DRDACKS at 244. The assertion at 244 causes the G multiplexer to again select the contents of the RPX pointer storage register for output on the address bus 40, thereby supplying the one currently stored as the RPX pointer to the input of the incrementer. The incrementer increments the pointer address one to pointer address two, and applies it to the output line 83. Upon completion of the DPC read service request at time t14, as shown at 246, the control logic causes the multiplexer in Figure 4 to select its number 1 input coupled to the output of the incrementer, thereby loading the address two pointer value into the RPX pointer storage register. This completes the illustration of the nonlinked-list transmission of words from the buffer memory. This example extends from time t0 to t14 in Figure 13.

Linked-List Transmission Example

At time t14 the data path controller asserts another read request as shown at 248. This is acknowledged by the control logic 125 by asserting the signal DRDACKS at time t15, as shown at 250. As described earlier, the assertion of the DRDACKS signal at time t15 causes the current contents of the RPX pointer storage register, address two, to be applied to the input of the incrementer and incremented to a value of address three. To best understand the operation of the RAM buffer controller in the linked-list mode, reference is made to Figure 14, which shows a hypothetical example of several memory locations having addresses 0-5 and showing the contents of these memory location. Figure 13 shows the data on the data bus during read cycles on the line entitled "Memory Contents on DBI". The first read transfer from the buffer memory 38 occurs between times t8 and t10. At that time the RPX pointer address is zero, and the contents of the zero address, a 100 value, is placed on the DBI bus. The next read cycle occurs between times t12 and t14. At that time, the RPX pointer value is one, and the contents of address one are placed on the DBI bus. Thus, a 200 value is found on the DBI bus between times t12 and t14. The next read cycle occurs between times t15 and t17. At that time the RPX pointer is on address location two, so the data on the DBI bus represents a value 300. The deactivation of the signal DRDACKS at time t17 causes the C multiplexer in Figure 4 to select its number 1 input, thereby loading the output of the incrementer, a value of three, into the RPX pointer storage register, as shown at 252 in Figure 13. At time t17, the data path controller again asserts a read request, as shown at 254. This read request is acknowledged at 256 and time t18. As described before, this causes the RPX pointer to be selected for output by the G multiplexer on the address bus 40, thereby outputting an address of three to the buffer memory. This address value of three is incremented by the incrementer 82 and applied to the first input of the C multiplexer by the line 83.

Assume now that the contents of address location number three are a pointer to the starting address of a packet number 2. Assume also that address locations zero through two contain data words from a packet number 1. Address location three will contain a unique delimiter that indicates to the data path controller that its contents are a pointer address. When the data path controller detects this delimiter (not shown in Figure 14), the data path controller asserts a signal LDRPXS at time t18, as shown at 258. This causes the current contents of the DBI bus to be loaded into the RPX register. As seen from Figure 2, there are two signals of this nature, LDRPXS and LDRPXA. These two signals are used to control the loading of the RPXS and RPXA pointer storage registers in Figure 5, respectively. Thus, upon deactivation of the signal DRDACKS at 260 and deactivation of the signal LDRPXS at 262, the control logic 125 causes the C multiplexer to switch such that the value on the DBI bus, a 100 is loaded into the RPX pointer storage register.

At time t20, the DPC asserts another read request, as shown at 262. This read request is acknowledged by assertion of the signal DRDACKS at time t21 as shown at 264. The acknowledgement of the DPC read request causes the control logic to send a select signal to the G multiplexer causing it to select the RPX contents for output on the address bus 40. Thus, the current contents of the RPX pointer storage register, i.e., an address 1000, is output on the address bus 40, and the contents of the memory location 1000 are then output on the DBI bus. This is shown in Figure 13 as the existence of a 20, the contents of memory location 1000, on the DBI bus between times t20 and t22.

The contents of memory location 1000 are not a pointer, but are the first word of packet number 2. Therefore, the data path controller does not activate the signal LDRPXS. The acknowledgement of the DPC read request at time t21 also applies the current contents of RPX to the incrementer input, thereby causing the incrementer to increment the 1000 address pointer to an address pointer having a value of 1001. Upon completion of the service of the DPC read request at time t23, the signal DRDACKS is deactivated, as shown at 266. This causes the C multiplexer to select its number 1 input, thereby loading the output of the incrementer, the value 1001, into the RPX pointer storage register. Processing or read requests then continue as described above. This completes the description of linked-list transmission protocol.

Another feature of the invention which is useful is pipelined arbitration. Referring to Figure 15 there is shown a conceptual diagram of the pipelined arbitration scheme of the invention. The basic idea is to overlap the arbitration cycle with the second service cycle of each two cycle transfer just preceding the service cycle. An example will clarify this. Assume at time t1, a read request comes in from the host system and from the node processor simultaneously. The arbitration logic decides the priority and selects the node processor read request for service starting at time t2 and extending for two cycles to time t4. The RPX pointer is selected, and its address remains constant for two cycles from t2 to t4. Meanwhile the host system read request remains pending, and a host system write request is asserted at time t2. At time t3, the arbitration logic will be caused by the control logic 125 to arbitrate between the host system read request

and the host system write request. This arbitration will occur between times t3 and t4. The winning request will be decided and ready at time t4 which is the same time that the service for the winning request decided between times t1 and t2 ends. Thus service of the winning request from round two begins immediately upon completion of service of the winning request from round 1.

MULTI DATA PATH CONTROLLER SYSTEMS

There are some applications where it is useful to allow two systems with a need to access a RAM and which may or may not be running synchronously with each other to share a single RAM buffer controller. Where a single clock drives all the systems and is coupled to the RAM buffer controller, there is no problem. Such a system is illustrated in Figure 1 where the single clock is derived from the incoming data stream on the medium 20 and the clock signal is applied to all units in the system so that all are running synchronously.

However, there are some applications where the two systems needing DMA access to the buffer memory whose address port is controlled by a single RAM buffer controller each have their own clock. Such an application is shown in Figure 16. In Figure 16, system 1 can be any system with DMA needs and its own clock, CLOCK 1, such as a host system, node processor, computer peripheral etc., and system 2 can be any other system with DMA needs and its own clock, CLOCK 2. The RAM buffer controller is coupled to only CLOCK 1, and it controls the address ports and supplied address signals and control signals, including CLOCK 1, to a buffer memory. System 1 is coupled to the data ports of the buffer memory by the data bus DATA 1. System 2 is coupled to the data ports of the buffer memory via its own data bus DATA 2, a reclocking latch and data bus DATA 1. CLOCK 2 is applied to the clock input of the reclocking latch such that when system 2 wishes to load data into or receive data from the buffer memory, the data is loaded in the reclocking latch in synchronization with the system 2 clock, even though the buffer memory is operating in synchronization with the system 1 clock. The purpose of the reclocking latch is to eliminate the timing difference between the DMA operations of systems 1 and 2.

Although the invention has been described in terms of the preferred embodiment and other embodiments described herein, those skilled in the art will appreciate many variations on the circuitry described herein which may be successfully used without departing from the true spirit and scope of the claims appended hereto. For example, the circuitry of Figures 4 and 5 could be replaced with a general purpose computer which is programmed to carry out the protocols described herein in Figure 13 and Figures 8-12. All such modifications are intended to be included within the scope of the claims appended hereto.

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Claims

1. An apparatus for generating pointer addresses for guiding data operations with a memory comprising:
 - first means for storing a plurality of pointer addresses;
 - second means coupled to said first means for selecting one of said pointer addresses for output to said memory;
 - third means coupled to said second means and to said first means for incrementing the selected pointer address and loading it back into said first means to replace the selected pointer address.
2. The apparatus of claim 1 wherein said first means stores a write pointer and read pointer for receive packets to be stored in said memory and a read pointer and a write pointer for transmit packets to be retrieved from said memory for transmission.
3. The apparatus of claim 2 further comprising an initialization bus coupled to said first means for allowing any of said pointers to be set at a specific value and fourth means coupled to said initialization bus and said first means for receiving and storing starting address pointers for receive and transmit FIFO memory organization in said memory and for causing said starting addresses to be stored in said first means at predetermined times.
4. The apparatus of claim 3 further comprising fifth means coupled to said initialization bus for receiving and storing end addresses for said receive and transmit FIFO memory organizations in said memory and sixth means coupled to said fifth means and to said second means for comparing the selected pointer address to the appropriate one of said end addresses and for generating an end signal indicating whether the selected pointer address is equal or not equal to the appropriate one of said end addresses.

5. An apparatus as defined in claim 4 further comprising a multiplexer having its output coupled to a data input of said first means and having an input coupled to an output of said fourth means for storing said starting address pointers and an input coupled to said initialization bus and having a control input and further comprising control means coupled to said control input and to said sixth means for receiving said end signal and for determining when the selected read pointer or write pointer has become equal to said end address and for causing said multiplexer to select the output of said fourth means such that the proper starting address pointer for the FIFO corresponding to the selected read or write pointer may be stored in said first means as the new value for the selected read or write pointer address.

6. An apparatus as defined in claim 5 further comprising full/empty comparison means coupled to said first means for comparing said read and write pointers for receive packets and generating a first signal indicating whether said pointers are equal and for comparing said read and write pointers for said transmit packets and generating a second signal indicating whether said pointers are equal.

7. An apparatus as defined in claim 6 further comprising means coupled to said full/empty comparison means for receiving said first and said second signals and for comparing these signals to the type of memory access that occurred in the memory access cycle just previous to the cycle that caused the compared pointers to become equal and to the type of memory access that occurred in the memory access cycle that caused the pointers to become equal and for making a determination in the case of said first signal indicating the status of the receive FIFO whether that FIFO is full or empty and for making a determination in the case of said second signal indicating the status of the transmit FIFO whether that FIFO is full or empty.

8. An apparatus as defined in claim 1 further comprising means coupled to said first means for initializing separate write and read pointers for a first stack and for initializing separate write and read pointers for a second stack and for controlling said third means to manage the incrementing and decrementing of said write and read pointers such that a last in, first out memory access operation in said memory may be performed.

9. An apparatus as defined in claim 4 wherein said fifth means includes means for allowing the end pointer addresses to be initialized or changed to any desired address within the address space of said memory.

10. An apparatus as defined in claim 2 wherein said first means includes a plurality of registers each for storing one of said pointer addresses one of which is for storing the read address for transmit packets and further comprising linked list means for allowing the read pointer for transmit packets to be changed to a new arbitrary address within the address space of said memory during each memory access cycle for purposes of allowing linked lists of packets to be transmitted.

11. An apparatus as defined in claim 10 further comprising an initialization bus for carrying the initial address for read pointer for transmit packets and a DBI bus for carrying linking addresses for storage in the register storing the read pointer address for transmit packets and wherein said linked list means includes an input multiplexer having its output coupled to the data input of said register for storing the read address for transmit packets and having an input coupled to initialization bus and having an input coupled to said DBI bus and having an input coupled to said third means and having a control input, and further comprising control means coupled to said control input for causing said multiplexer to select the input coupled to said DBI bus during transmission of a packet organized as a linked list in said memory such that linking addresses from said DBI bus will be loaded in the register storing the read pointer for transmit packets.

12. An apparatus as defined in claim 10 wherein said first means includes at least two registers for storing separate and individual read address pointers for transmit packets organized into at least two linked lists and wherein said linked list means includes means for allowing separate linking addresses for said at least two linked lists to be stored in each of said at least two registers where each register and linking address correspond to one of said at least two linked lists.

13. An apparatus as defined in claim 11 wherein said first means includes at least two registers for storing separate and individual read address pointers for transmit packets organized into at least two linked lists and wherein said linked list means includes means for allowing separate linking addresses for said at least two linked lists to be stored in each of said at least two registers where each register and linking address correspond to one of said at least two linked lists.

14. An apparatus as defined in claim 1 wherein said first means is comprised of a plurality of registers for storing said pointers and further comprising a plurality of multiplexers each having an output coupled to a data input of one of said registers and each having a plurality of inputs, and further comprising an initialization bus coupled to selected ones of said inputs and wherein said third means is coupled to other selected ones of said inputs.

15. An apparatus as defined in claim 14 wherein one said registers is a random access memory address register.

16. An apparatus as defined in claim 5 further comprising a random access memory address register coupled to said initialization bus and wherein said control means ignores said end signal when said second means selects the output of said random access memory address register for output to said memory.

17. An apparatus as defined in claim 1 wherein said first means includes a plurality of registers for storing said pointer addresses including a first register for storing a write pointer for receive packets and having an output coupled to said second means and a second register for storing a shadow write pointer for receive packets where said first and second registers have inputs and outputs, and further comprising first and second multiplexers having outputs coupled respectively to said inputs of said first and second registers, each said multiplexer having a plurality of inputs wherein one input of said first multiplexer is coupled to said third means for receiving the incremented value of said write pointer for receive packets when said second means selects the write pointer for receive packets for output to said memory and wherein the output of said second register is coupled to an input of each of said first and second multiplexers and wherein the output of said first register is coupled to an input of each of said first and second multiplexers.

18. An apparatus as defined in claim 17 wherein said first and second multiplexers also have control inputs and further comprising control means coupled to said control inputs and having a first input for receiving a FLSPKT signal indicating a received packet or portion thereof is to be flushed and having a second input for receiving a BRCVPKT signal indicating packet reception is beginning and having a third input for receiving a DWRREQ signal indicating that packet reception is occurring, said control means for causing said second multiplexer to select the input coupled to the output of said first register and for causing said first multiplexer to select the input coupled to said third means to cause said first register to load the write pointer for receive packets as incremented by said third means upon receiving said BRCVPKT signal and for causing said first multiplexer to continue to select said input coupled to said third means upon receiving said DWRREQ signal and for causing said first multiplexer to select the input coupled to the output of said second register and thereafter to cause said first multiplexer to select said input coupled to said third means upon receiving said FLSPKT signal.

19. An apparatus as defined in claim 17 wherein said first and second multiplexers also have control inputs and further comprising control means coupled to said control inputs, and having a first input for receiving a FLSPKT signal indicating a received packet or portion thereof is to be flushed, and having a second input for receiving a BRCVPKT signal indicating packet reception is beginning, and having a third input for receiving a DWRREQ signal indicating that packet reception is occurring, and having a fourth input for receiving an ERCVPKT signal indicating a packet has been successfully received, said control means for causing said second multiplexer to select the input coupled to the output of said first register to cause the then existing write pointer for receive packets to be saved in the second register and for causing said first multiplexer to select the input coupled to said third means to cause said first register to load the write pointer for receive packets as incremented by said third means upon receiving said BRCVPKT signal and for causing said first multiplexer to continue to select said input coupled to said third means upon receiving said DWRREQ signal such that the write pointer for receive packets is incremented on every clock cycle during packet reception and, upon receiving said ERCVPKT signal, for causing said first multiplexer to select the input coupled to the output of said second register and for causing said second multiplexer to select the input coupled to the output of said first register such that the contents of these two register are exchanged, and for causing said second means to select the output of said first register for application to the address inputs of said memory, and further comprising data path controller means coupled to the data inputs of said memory for keeping track of the length of the received packet and its status and for writing this data into the location in memory pointer to by the current contents of said first register, said control means also for causing said first multiplexer to select the input coupled to the output of said register after said status and length data is written so as to cause said first register to load the contents of said second register to restore the write pointer for receive packets to its state as it existed at the time of receiving said ERCVPKT signal, said data path controller for writing a predetermined bit in the memory location addressed by the contents of said first register and for subsequently causing said first multiplexer to select the input coupled to the output of said first register until the next BRCVPKT signal is received.

20. An apparatus as defined in claim 17 wherein said first and second multiplexers also have control inputs and further comprising control means coupled to said control inputs, and having a first input for receiving a FLSPKT signal indicating a received packet or portion thereof is to be flushed, and having a second input for receiving a BRCVPKT signal indicating packet reception is beginning, and having a third input for receiving a DWRREQ signal indicating that packet reception is occurring, and having a fourth input

for receiving an RCVABRT signal indicating packet that has been received or partially received is to be aborted but its status and length information is to be recorded, said control means for causing said second multiplexer to select the input coupled to the output of said first register to cause the then existing write pointer for receive packets to be saved in the second register and for causing said first multiplexer to select the input coupled to said third means to cause said first register to load the write pointer for receive packets as incremented by said third means upon receiving said BRCVPKT signal, and for causing said first multiplexer to continue to select said input coupled to said third means upon receiving said DWRREQ signal such that the write pointer for receive packets is incremented on every clock cycle during packet reception and, upon receiving said RCVABRT signal, for causing said first multiplexer to select the input coupled to the output of said second register such that the contents of the second register is loaded in said first register, and for causing said second means to select the output of said first register for application to the address inputs of said memory, and further comprising data path controllers means coupled to the data inputs of said memory for keeping track of the length of the received packet and its status and for writing this data into the location in memory pointed to by the current contents of said first register, said control means also for causing said first multiplexer to select the input coupled to the output of said third means after said status and length data is written so as to cause said first register to reload the current write pointer for receive packets as incremented by said third means, said data path controller for writing a predetermined bit in the memory location addressed by the contents of said first register after its contents have been incremented and for subsequently causing said first multiplexer to select the input coupled to the output of said first register until the next BRCVPKT signal is received.

21. An apparatus as defined in claim 11 further comprising a starting address register for storing a new link address pointing to the location in memory of a predetermined packet and having an input coupled to said initialization bus and having an output coupled at least to an input of the source multiplexer having its output coupled to the input of the register storing the read pointer for the transmit packets, said source multiplexer having a control input, and further comprising a control means coupled to said control input and having a control input for receiving a signal INICLBN indicating that said predetermined packet should be sent, said control means for causing said source multiplexer to select the input coupled to the output of said starting address register when said INICLBN signal is received.

22. An apparatus as defined in claim 21 further comprising an end address register having an input and a second source multiplexer having an input coupled to the output of said register for storing the read address pointer for the transmit packets and having its output coupled to said input of said end address register and having a control input, and wherein said control means is also coupled to said control input of said second source multiplexer and wherein said control means is also for simultaneously causing said second source multiplexer to select said input coupled to the output of said register for storing the read address pointer for transmit packets simultaneously with the selection by said source multiplexer of its input coupled to the input of said starting address register upon receiving said INICLBN signal such that the link address is loaded in said register which is for storing the read pointer for transmit packets simultaneously with the loading of the current read pointer for transmit packets into the end address register.

23. An apparatus as defined in claim 1 further comprising arbitration means coupled to said second means and having a plurality of service request signal inputs for receiving signals indicating that a particular pointer is desired to be selected, said arbitration means for arbitrating said request signals and selecting the currently asserted request with the highest priority.

24. An apparatus as defined in claim 23 further comprising a clock input for receiving a clock signal and wherein said second means has a control input and wherein said first means comprises a plurality of registers for storing predetermined ones of said pointer addresses, each register having an input and an output and a clock input for receiving said clock signal for causing the data at said input to be latched into said register and to be available for output on said output, and further comprising a plurality of multiplexers each having its output coupled to a said input of one of said registers and each having at least one input coupled to the output of the register to which the multiplexer has its output coupled, and wherein said third means has an output carrying the incremented value of the pointer selected by said second means, said output of said third means coupled to another input of each said multiplexer, and each said multiplexer having a control input, and further comprising a control means coupled to said control inputs of said multiplexers and to said arbitration means and to a control input of said second means for sensing which request signal has received priority and for causing said second means to select the appropriate one of said pointers in accordance with the result of the arbitration and for causing said multiplexer coupled to the input of the register whose pointer contents have been selected for output by said second means to select its

input coupled to the output of said third means for a first clock cycle and for causing said multiplexer coupled to the selected register to select the multiplexer input coupled to the output of the selected register during a second clock cycle.

25. The apparatus of claim 24 wherein said control means also causes said arbitration means to perform its next arbitration for the next pointer to be selected during said second clock cycle during which the currently selected pointer is selected for output to the address inputs of said memory.

26. An apparatus as defined in claim 1 wherein said first means includes a plurality of registers for storing said pointers and further comprises an initialization bus coupled to an input of each register and wherein one of said registers is for storing a random access pointer address which can be set initially by data on said initialization bus and will be incremented by said third means when said random access pointer is selected for output by said second means.

27. An apparatus as defined in claim 26 further comprising arbitration means having a plurality of inputs for coupling to conductors carrying request signals and coupled to said second means for arbitrating any simultaneously asserted requests and selecting the highest priority request and for causing said second means to select said random access pointer register output for output to said memory.

28. An apparatus as defined in claim 26 further comprising instruction decoder means having an instruction bus input and coupled to said second means for decoding an instruction received at said instruction bus input and causing said random access pointer register output to be selected for output.

29. An apparatus as defined in claim 27 further comprising means having an instruction bus input and coupled to said second means for decoding an instruction received at said instruction bus input and causing said random access pointer register output to be selected for output.

30. An apparatus as defined in claim 28 wherein said instruction decoder means is also coupled to said third means in such a way that when a predetermined instruction is received, the third means can be stopped from incrementing the selected pointer.

31. An apparatus as defined in claim 1 further comprising an initialization bus and wherein said first means is comprised of a plurality of registers each having an input and further comprising a plurality of multiplexers having their outputs coupled to said register inputs, each multiplexer having an input coupled to said third means for receiving the incremented value of the selected pointer and each multiplexer having an input coupled to said initialization bus.

32. An apparatus as defined in claim 1 further comprising an initialization bus and wherein said first means is comprised of a plurality of registers each having an input and further comprising a plurality of input multiplexers having their outputs coupled to said register inputs, and further comprising a source multiplexer having an input coupled to said initialization bus and having an input coupled to said third means and having an output, each said input multiplexer having an input coupled to said output of said source multiplexer.

33. A system comprising:

a clock;

a first system for reading and writing data coupled to said clock and having a first data bus and having a first output line for carrying a first memory access request signal;

40 a second system for reading and writing data coupled to said clock and having a second data bus and having a second output line for carrying a second memory access request signal;

a memory coupled to said clock and having a data port and an address port, said data port coupled to said first and second data buses; and

45 a buffer controller means coupled to said clock and coupled to said address port and to said first and second memory access request signals for arbitrating said memory access requests and for generating pointer addresses at said address port for controlling access to data stored in said memory via said data buses.

34. An apparatus as defined in claim 33 wherein said buffer controller means is for managing said pointer addresses such that said memory appears to each of said first and second systems as a pair of independently operating FIFO memories.

35. A system comprising:

a first system for reading and writing data having a first clock and having a first data bus and having a first output line for carrying a first memory access request signal;

55 a second system for reading and writing data having a second clock and having a second data bus and having a second output line for carrying a second memory access request signal;

a memory coupled to said first clock and having a data port and an address port, said data port coupled to said first data bus;

a buffer controller means coupled to said first clock and coupled to said address port and to said first and

second memory access request signals for arbitrating said memory access requests and for generating pointer addresses at said address port for controlling access to data stored in said memory via said data buses and having an enable output, said buffer controller means also for generating an enable signal at said enable output indicating when a pointer for said second system has been output to said memory port; and
 5 a latch having a data input coupled to said second data bus and having a data output coupled to said data port and having a clock input coupled to said second clock and having an output enable input coupled to said enable output of said buffer controller means for receiving said output enable signal to cause data stored in said latch by said second system to be applied to said data port of said memory.

36. An apparatus as defined in claim 35 wherein said buffer controller means is for managing said
 10 pointer address such that said memory appears to each of said first and second systems as a pair of independently operating FIFO memories.

37. An apparatus for generating pointer addresses at an output for coupling to an address port of a memory for controlling data operations of said memory comprising:
 a clock means for generating a periodic series of load signals;
 15 a first register having an input and having a load input coupled to said clock for causing said register to load and store the address at its input upon each occurrence of said load signal and having an output, said register for storing a read pointer for a transmit FIFO;
 a second register having an input and having a load input coupled to said clock for causing said register to load and store the address at its input upon each occurrence of said load signal and having an output, said
 20 register for storing a write pointer for a transmit FIFO;
 a third register having an input and having a load input coupled to said clock for causing said register to load and store the address at its input upon each occurrence of said load signal and having an output, said register for storing a read pointer for a receive FIFO;
 25 a fourth register having an input and having a load input coupled to said clock for causing said register to load and store the address at its input upon each occurrence of said load signal and having an output, said register for storing a write pointer for a receive FIFO;
 an output multiplexer having inputs coupled to said outputs of said first, second, third and fourth registers and having an output and having a control input for receiving a signal controlling which of said inputs is coupled to said output;
 30 an input multiplexer for each of said first through fourth registers, each input multiplexer having an output coupled to said input of its associated register and each input multiplexer having first and second inputs and having a control input for receiving a signal controlling which of said inputs is coupled to said outputs;
 an incrementing means having an input coupled to the output of said output multiplexer and an output coupled to each first input of each said input multiplexer for incrementing the pointer address stored in the
 35 selected register and supplying the incremented result to the inputs of said input multiplexers;
 an initialization bus coupled to each second input of said input multiplexers for allowing any of said registers to be loaded with an address placed on said bus; and
 control means having an instruction port and output signal lines coupled to said control inputs for said input multiplexers to cause any selected one to select its second input for coupling to said output of said input
 40 multiplexer when a predetermined instruction arrives at said instruction port;
 arbitration means having a plurality of memory access request inputs and having a plurality of control signal lines coupled to said control inputs of said input multiplexers and to said control input of said output multiplexer for arbitrating simultaneous memory access requests, selecting the highest priority current request and activating predetermined signal lines to cause said output multiplexer to select the output of the
 45 register storing the pointer selected as the current highest priority by said arbitration means for output to said memory and for causing the input multiplexer coupled to the selected register to select the input coupled to said incrementing means such the current pointer in the selected register will be incremented and the new value will be loaded on the next clock cycle.

38. An apparatus as defined in claim 37 further comprising means coupled to said initialization bus and
 50 to said output of said output multiplexer and having a control output for storing an end address for each of said transmit FIFO and said receive FIFO and for comparing the address of the currently selected FIFO to the appropriate one of the end addresses and for generating a signal indicating whether or not the end address has been equalled.

39. An apparatus as defined in claim 38 further comprising random access memory means coupled to
 55 said output multiplexer and said incrementing means and to said initialization bus and to said control means for allowing any address on said initialization bus to be stored in said random access memory means and

output to said memory upon receipt of a predetermined instruction by said control means and to allow the address from the initialization bus to be incremented or not incremented and re-output depending upon the instruction received.

40. The apparatus of claim 37 further comprising a data bus and a means coupled to said output
5 multiplexer and selectively coupled to said data bus for storing a second read pointer for a transmit FIFO.

41. The apparatus of claim 37 further comprising start address registers selectively coupled to said initialization bus for storing start addresses for said transmit and said receive FIFO's and having an output coupled to inputs of said input multiplexers coupled to said first through fourth input multiplexers.

42. The apparatus of claim 37 further comprising a shadow write pointer register having a load input
10 coupled to said clock and having an input which is selectively coupled to the output of said fourth register through a fifth multiplexer having a control input coupled to said control means and having an output which is selectively coupled to the input of said register through said fourth multiplexer and wherein said control means also generates signals at a predetermined time to cause the fourth and fifth multiplexers to switch so as to cause said fourth register and said shadow write pointer register to swap their contents at
15 predetermined times.

43. A method of operating a RAM buffer controller for generating addresses for a memory, said controller including a first register, a second register, an incrementer and an output gate for gating the output of the first register to the address inputs of said memory comprising the steps of:
storing an address in said first register;

20 upon receipt of a first signal, storing the contents of said first register in said second register and incrementing the contents of said first register;

upon receipt of a second signal, holding the contents of said second register constant while incrementing the contents of said first register once for every access to said memory;

25 upon receipt of a third signal, simultaneously causing the contents of said first and second registers to be exchanged;

gating the output of said first register to said address inputs and writing data into the location in memory addressed by the contents of said first register;

causing the contents of said second register to be stored in said first register;

30 gating the output of said first register to said address ports of said memory and writing a predetermined bit in the memory location addressed by the contents of said first register with predetermined data;

holding the contents of said first register constant.

44. A method of operating a RAM buffer controller for generating addresses for a memory, said controller including a first register, a second register, an incrementer and an output gate for gating the output of the first register to the address inputs of said memory comprising the steps of:

35 storing an address in said first register;

upon receipt of a first signal, storing the contents of said first register in said second register and incrementing the contents of said first register;

upon receipt of a second signal, holding the contents of said second register constant while incrementing the contents of said first register once for every access to said memory;

40 upon receipt of a third signal, causing the contents of said second register to be stored in said first register.

45. A method of operating a RAM buffer controller for generating addresses for a memory, said controller including a first register, a second register, an incrementer and an output gate for gating the output of the first register to the address inputs of said memory comprising the steps of:

storing an address in said first register;

45 upon receipt of a first signal, storing the contents of said first register in said second register and incrementing the contents of said first register;

upon receipt of a second signal, holding the contents of said second register constant while incrementing the contents of said first register once for every access to said memory;

50 upon receipt of a third signal, causing the contents of said second register to be stored in said first register;

gating the output of said first register through to said address inputs and writing predetermined data in the location in memory addressed by the contents of said first register;

incrementing the contents of said first register;

gating the output of said first register through to said address inputs and writing a predetermined data bit into the memory location addressed by the contents of said first register; and

55 preventing the incrementing of the contents of said first register in preparation for the next memory access.

46. A method of operating a RAM buffer controller controlling the address inputs of a memory, said RAM buffer controller having a first register for storing write pointer addresses for use in storing received packets from a network in a memory and having a second register for storing a shadow write pointer and

having first and second multiplexers controlling the data which is input respectively to said first register and to said second register and having an incrementer having an input and having an output coupled to an input of said first multiplexer where said first multiplexer also has inputs coupled to the output of said first register and to the output of said second register and where said second register has inputs coupled to the output of said first register and to the output of said second register and having control logic with a plurality of input signal lines and a plurality of output control lines to control the operation of said first and second multiplexers where each of said first and second multiplexers has a clock input coupled to a clock generating load signals and having an output multiplexer with an input coupled to the output of said first register and an output coupled to the input of said incrementer comprising the steps of:

5 generating signals upon receipt of a first signal by said control logic to cause said second multiplexer to select its input coupled to the output of said first register and to cause the contents of said first register to be stored in said second register and thereafter generating signals to cause said second multiplexer to select its input coupled to the output of said second register and generating control signals to cause said first multiplexer to select its input coupled to said incrementer and thereafter generating signals causing

10 said first multiplexer to select its input coupled to the output of said first register; generating signals upon receipt of a second signal by said control logic to cause said output multiplexer to select the input coupled to the output of said first register once during every memory cycle and to cause said first multiplexer to again to select its input coupled to said incrementer on every memory access cycle and thereafter to hold the contents of said first register constant by selecting the input coupled to the output

15 of said first register between memory cycles, and generating signals to cause said second multiplexer to select its input coupled to the output of said second register while said first multiplexer is selecting its input coupled to said incrementer; generating signals upon receipt by said control logic of a third signal to cause said second multiplexer to select its input coupled to the output of said first register and to cause the contents of said first register to be stored in said second register and to cause the contents of said second register to thereafter be held

20 constant for a predetermined period and simultaneously with the loading of the contents of said first register into said second register to cause said first multiplexer to select its input coupled to the output of said second register and to cause the contents of said second register to be stored in said first register and thereafter to cause the contents of said first register to be held constant for a predetermined period;

25 generating signals to cause said output multiplexer to select the input coupled to the output of said first register and writing status and length information into the memory location in said memory addressed by the contents of said first register; generating signals causing said first multiplexer to select its input coupled to the output of said second register so that the contents of said second register will be loaded into said first register and to cause the

30 contents of said first register to be held constant; writing a predetermined bit into the memory location addressed by the contents of said first register; and generating signals causing said first multiplexer to select its input coupled to the output of said first register to hold the contents of said first register constant.

47. A method of operating a RAM buffer controller controlling the address inputs of a memory, said

40 RAM buffer controller having a first register for storing write pointer addresses for use in storing received packets from a network in a memory and having a second register for storing a shadow write pointer and having first and second multiplexers controlling the data which is input respectively to said first register and to said second register and having an incrementer having an input and having an output coupled to an input of said first multiplexer where said first multiplexer also has inputs coupled to the output of said first register

45 and to the output of said second register and where said second register has inputs coupled to the output of said first register and to the output of said second register and having control logic with a plurality of input signal lines and a plurality of output control lines to control the operation of said first and second multiplexers where each of said first and second multiplexers has a clock input coupled to a clock generating load signals and having an output multiplexer with an input coupled to the output of said first

50 register and an output coupled to the input of said incrementer comprising the steps of: generating signals upon receipt of a first signal by said control logic to cause said second multiplexer to select its input coupled to the output of said first register and to cause the contents of said first register to be stored in said second register and thereafter generating signals to cause said second multiplexer to select its input coupled to the output of said second register and generating control signals to cause said

55 first multiplexer to select its input coupled to said incrementer and thereafter generating signals causing said first multiplexer to select its input coupled to the output of said first register; generating signals upon receipt of a second signal by said control logic to cause said output multiplexer to select the input coupled to the output of said first register once during every memory cycle and to cause

said first multiplexer to again to select its input coupled to said incrementer on every memory access cycle and thereafter to hold the contents of said first register constant by selecting the input coupled to the output of said first register between memory cycles, and generating signals to cause said second multiplexer to select its input coupled to the output of said second register while said first multiplexer is selecting its input

5 coupled to said incrementer;

generating signals upon receipt by said control logic of a third signal to cause said first multiplexer to select its input coupled to the output of said second register and to cause the contents of said second register to be stored in said first register and thereafter to cause the contents of said first register to be held constant for a predetermined period.

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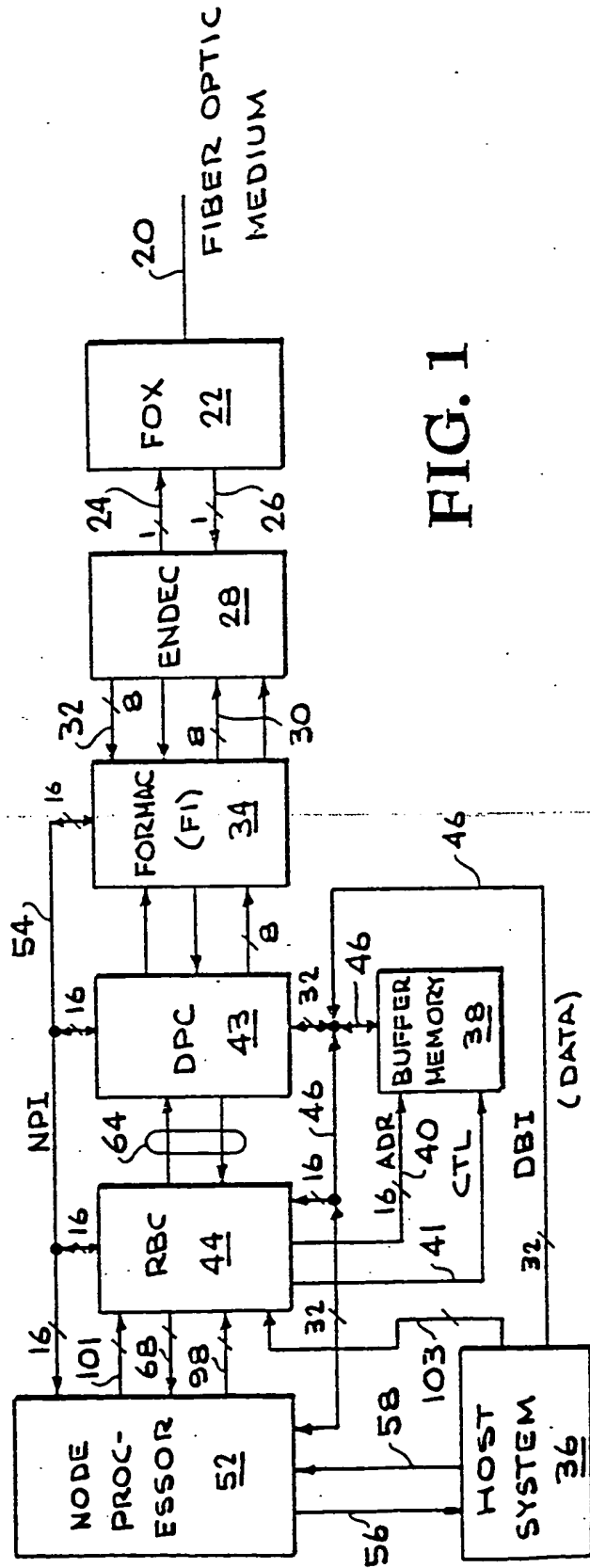


FIG. 1

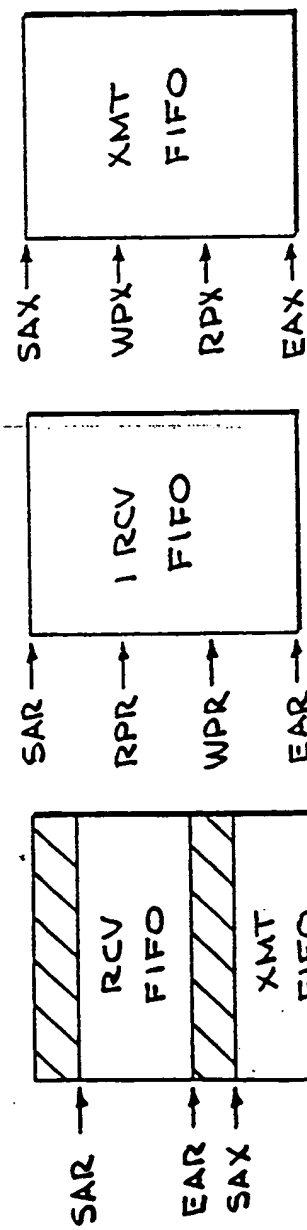


FIG. 3

FUNCTIONAL ORGANIZATION
OF RCV & XMT FIFO'S

PHYSICAL MEMORY
ORGANIZATION

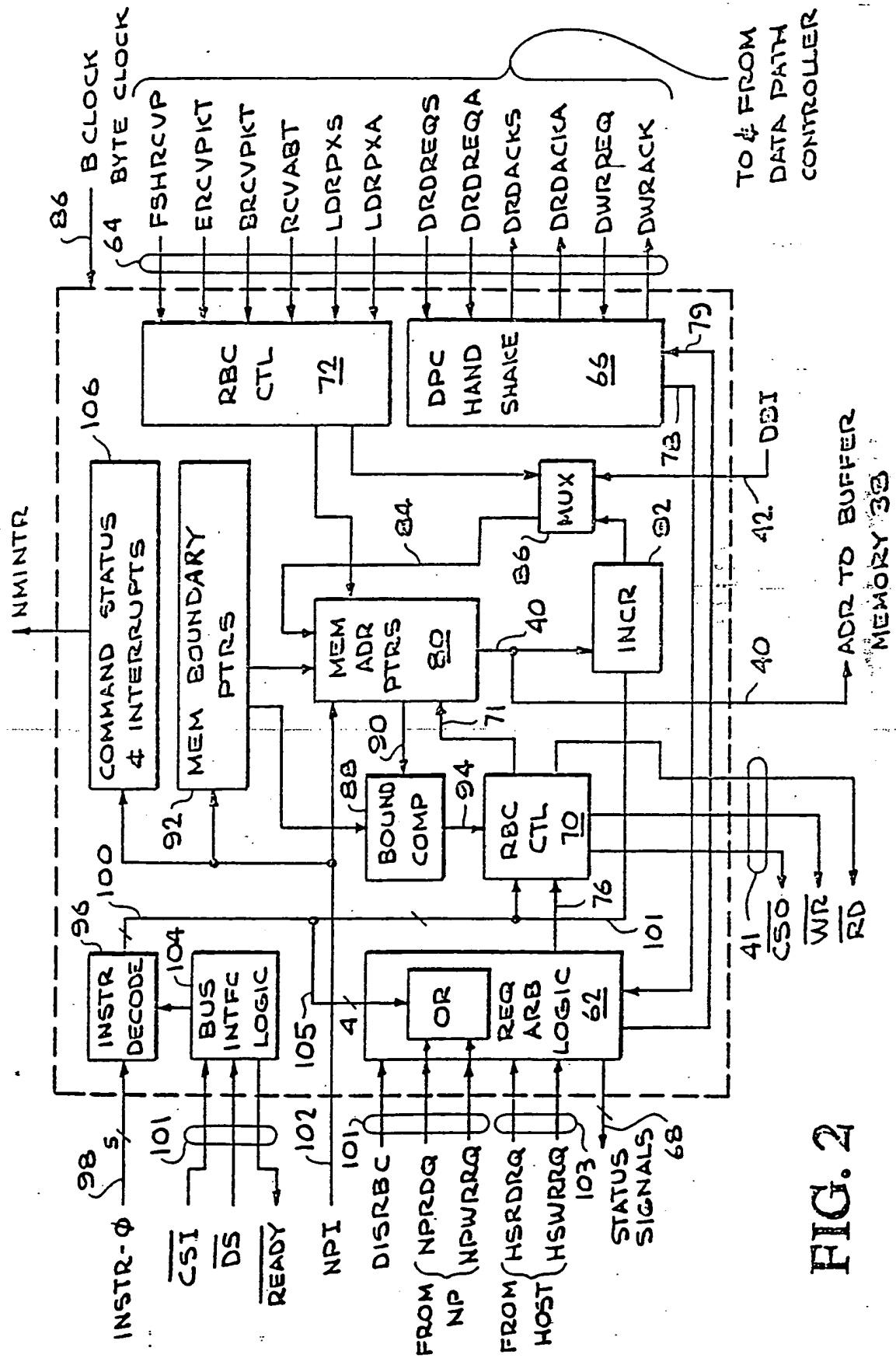


FIG. 2

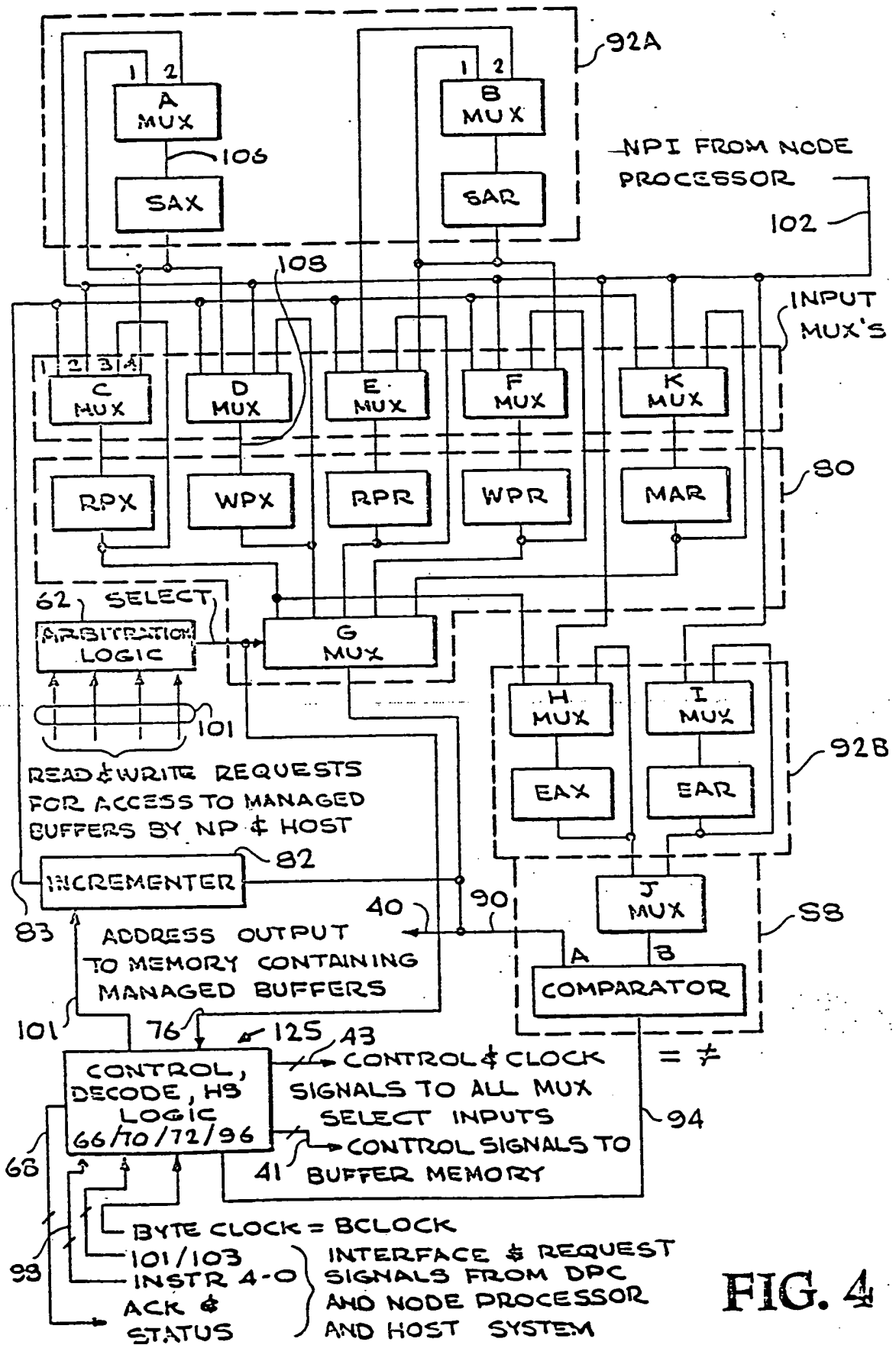
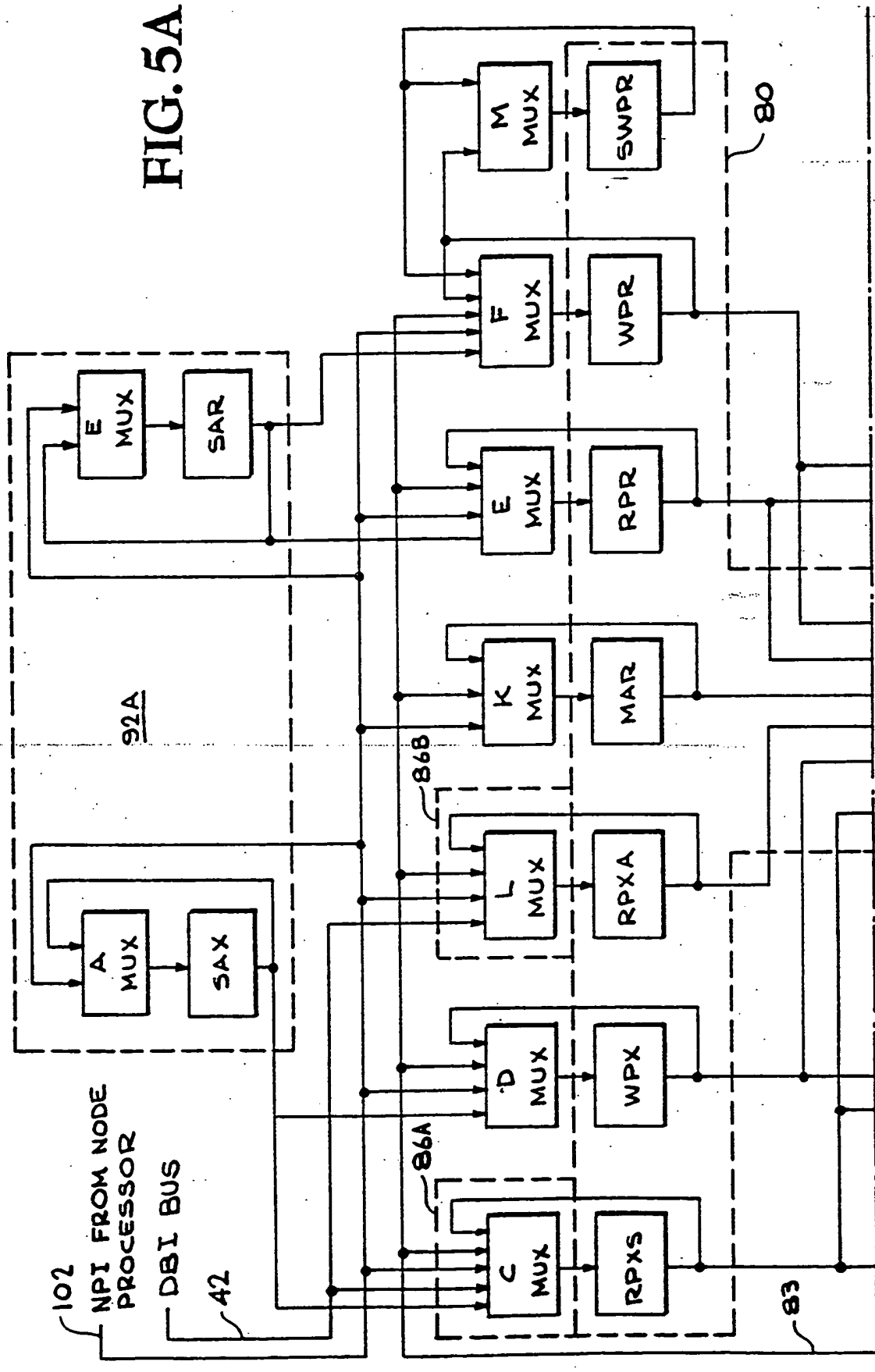
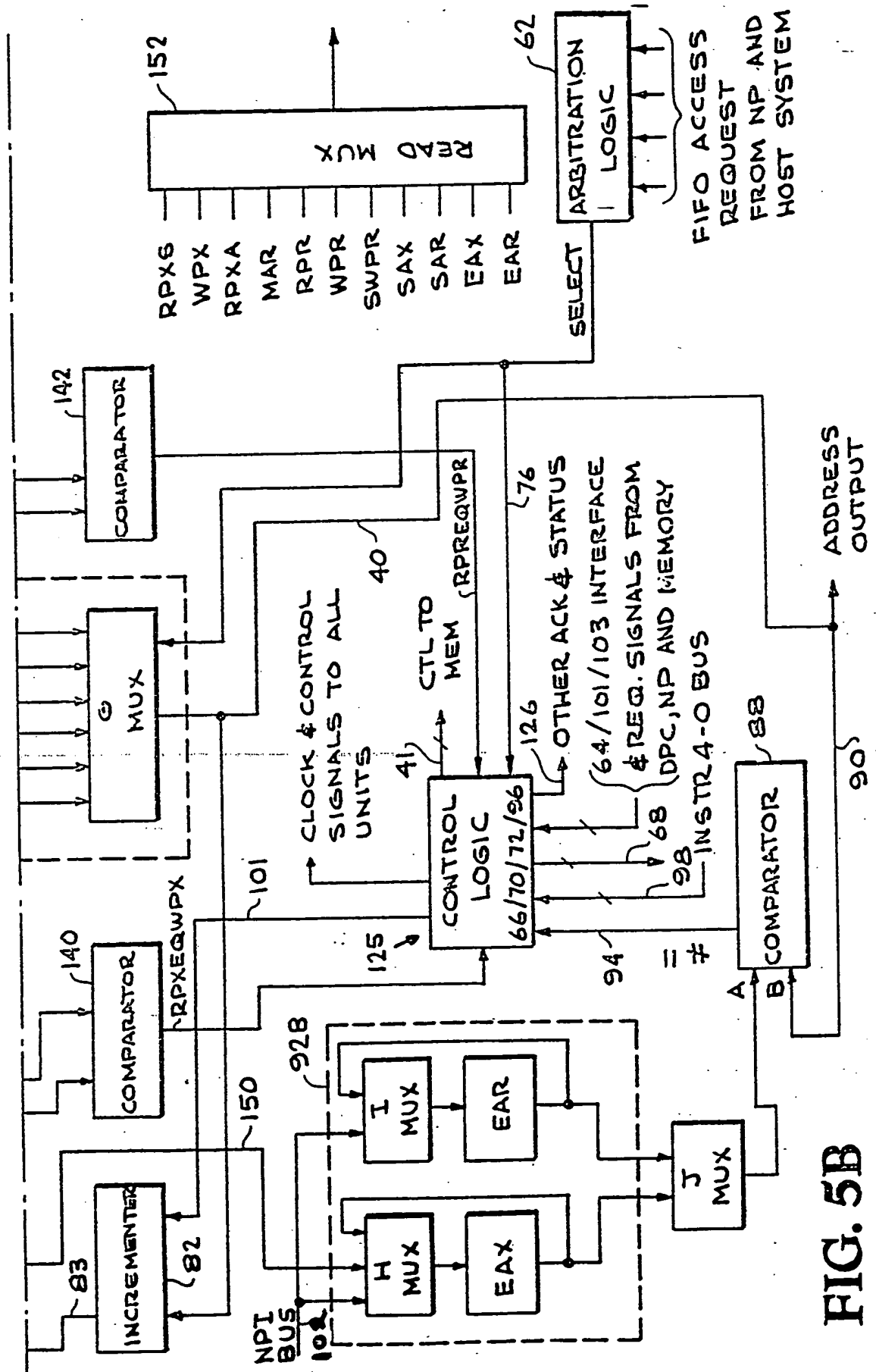


FIG. 4

FIG. 5A





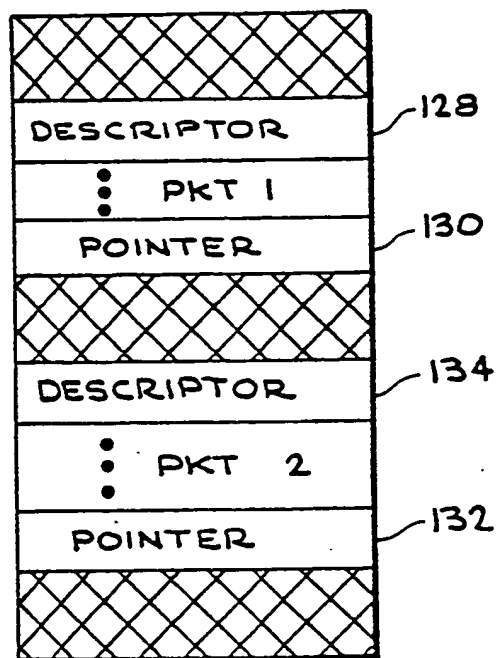


FIG. 6

CONTROL	DEST. ADR.	SOURCE ADR	DATA	CRC
---------	---------------	---------------	------------	-----

FIG. 7

	RECEIVE A NORMAL PACKET
1	AT START ASSUME: WPR=VALID ADR (INITIALIZE & HOLD) SWPR=UNKNOWN CONTENTS OF MEMORY LOCATION POINTED TO BY WPR = ZERO
2	ON ASSERTION OF BRCPKT: WPR → SWPR WPR+1 → WPR (SKIP 1ST LOCATION AFTER SAVING WPR)
3	DURING PACKET RECEPTION: DWRREQ IS ASSERTED AND WPR+1 → WPR SWPR → SWPR (HOLD)
4	ON ASSERTION OF ERCPKT: SWPR → WPR SIMULTANEOUSLY WITH WPR → SWPR
5	WRITE STATUS & LENGTH USING WPR
6	RESTORE WPR = SWPR → WPR
7	WRITE "ZERO" BIT USING WPR AND WPR → WPR (HOLD)

FIG. 8

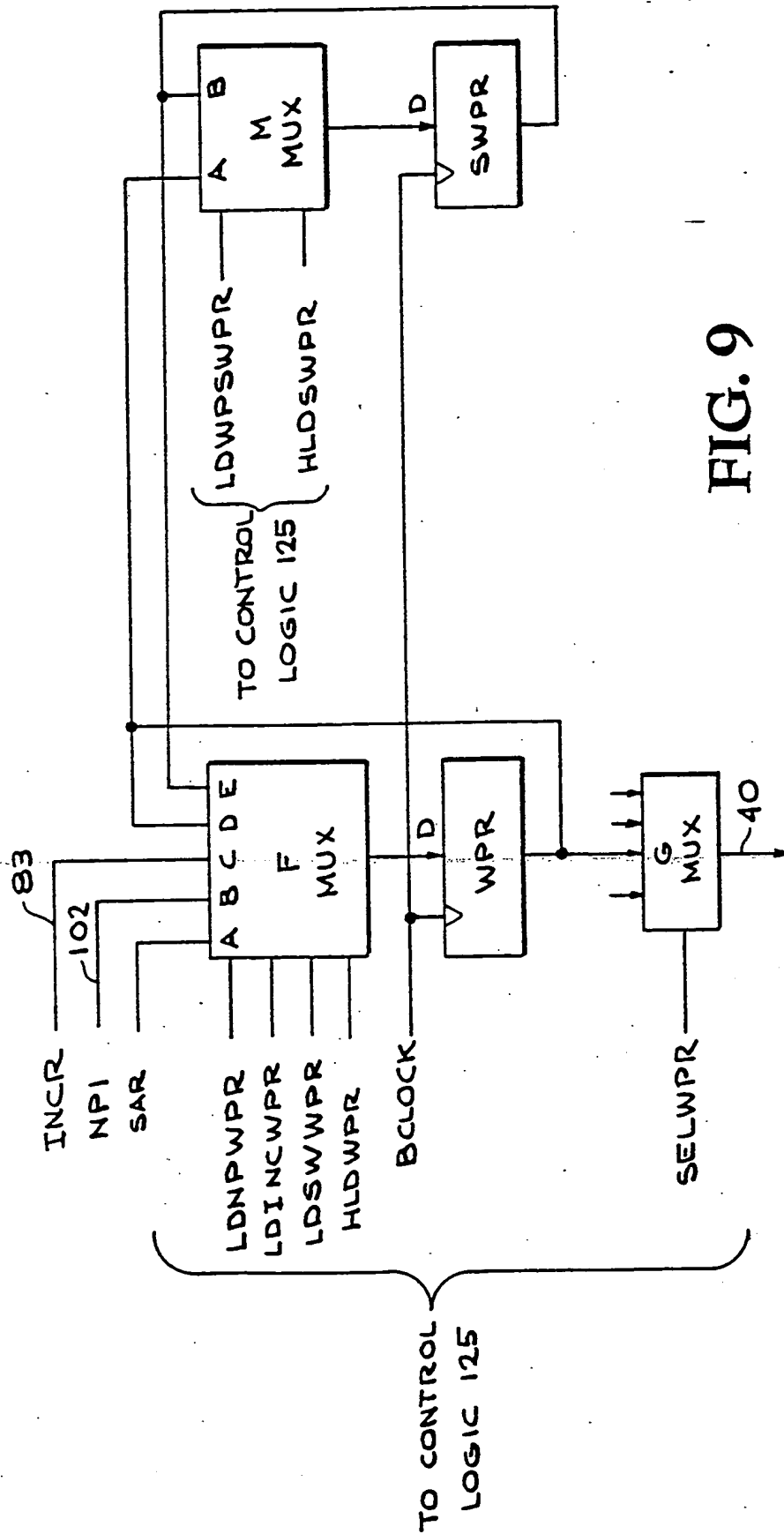
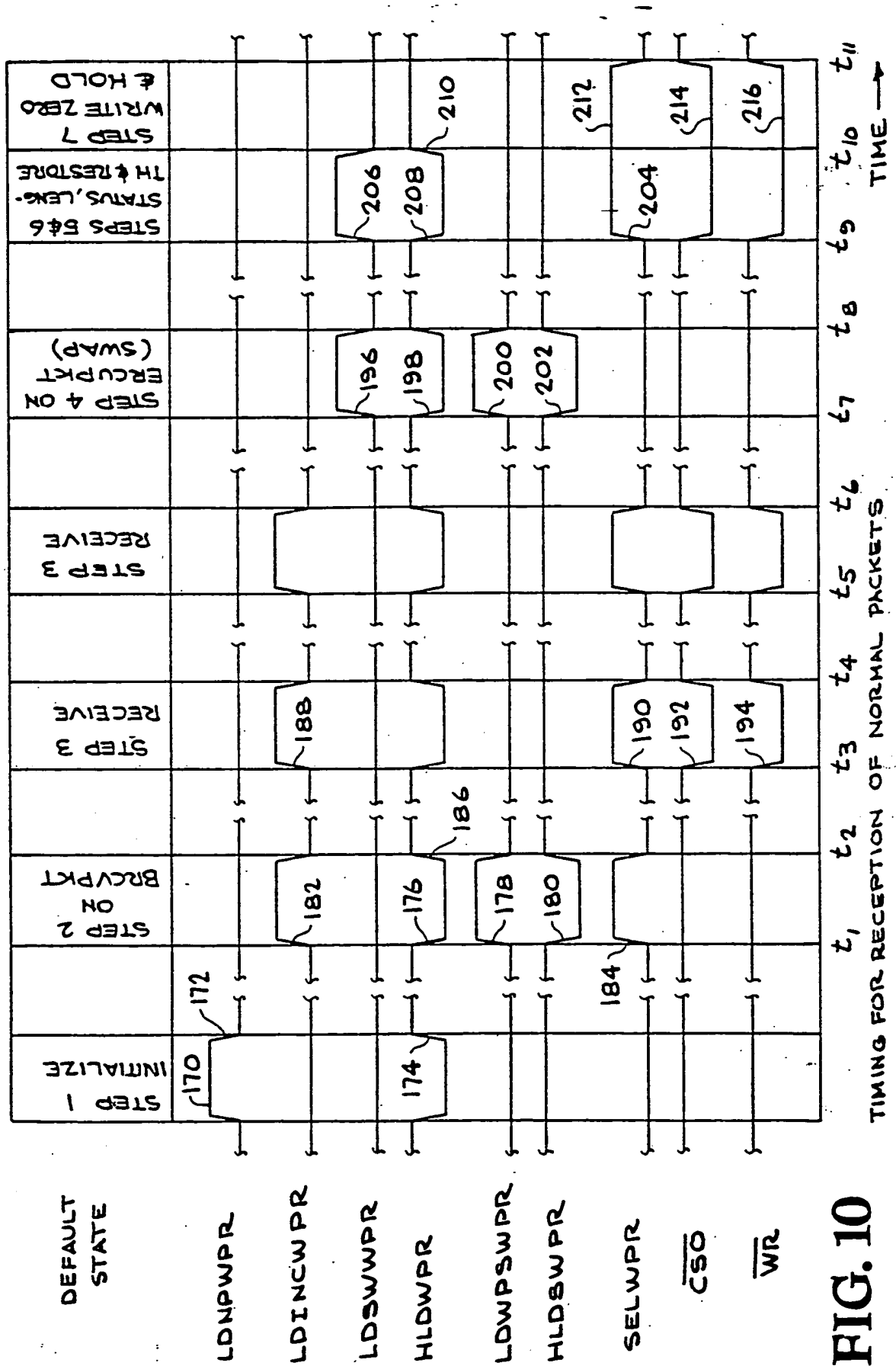


FIG. 9



FLUSHED PACKETS	
1	AT START ASSUME: WPR = VALID ADR SWPR = UNKNOWN CONTENTS OF MEMORY LOCATION POINTED TO BY WPR = ZERO
2	ON ASSERTION OF BRCVPKT: WPR → SWPR WPR + 1 → WPR
3	DURING PACKET RECEPTION: DWRREQ IS ASSERTED AND WPR + 1 → WPR SWPR → SWPR (HOLD)
4	ON ASSERTION OF FSHRCVP: SWPR → WPR

FIG. 11

ABORTED PACKETS	
1	AT START ASSUME: WPR = VALID ADR SWPR = UNKNOWN CONTENTS OF MEMORY LOCATION POINTED TO BY WPR = ZERO
2	ON ASSERTION OF BRCVPKT: WPR → SWPR WPR + 1 → WPR
3	DURING PACKET RECEPTION: DWRREQ IS ASSERTED AND WPR + 1 → WPR SWPR → SWPR
4	ON ASSERTION OF RCVABRT: SWPR → WPR
5	WRITE STATUS & LENGTH USING WPR & WPR + 1 → WPR
6	WRITE ZERO USING WPR AND WPR → WPR (HOLD)

FIG. 12

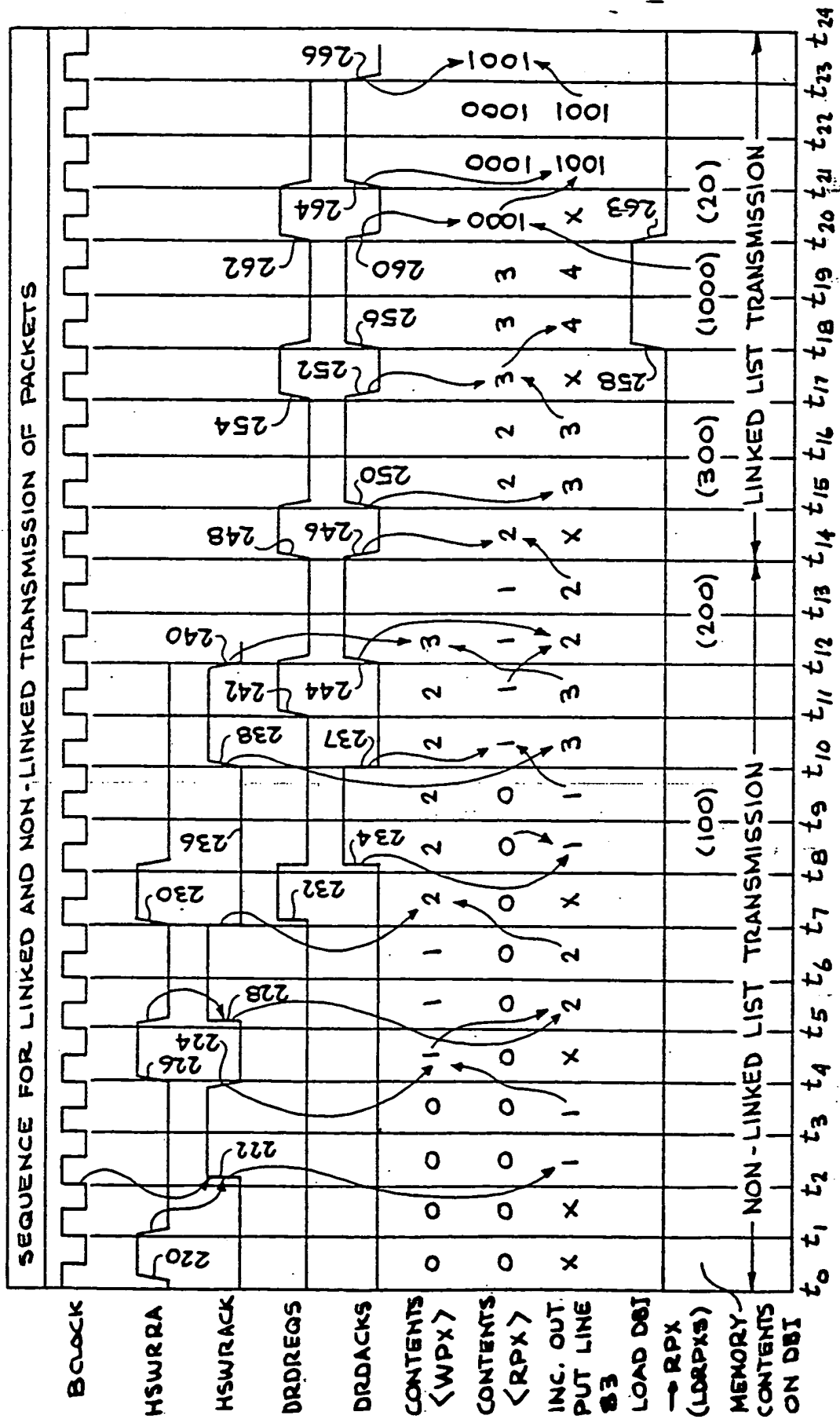


FIG. 13

0	100	} PACKET 1
1	200	
2	300	
3	1000	
4		
5		} POINTER TO STARTING ADDRESS OF PACKET 2
	⋮	
1000	20	
		} PACKET 2

FIG. 14

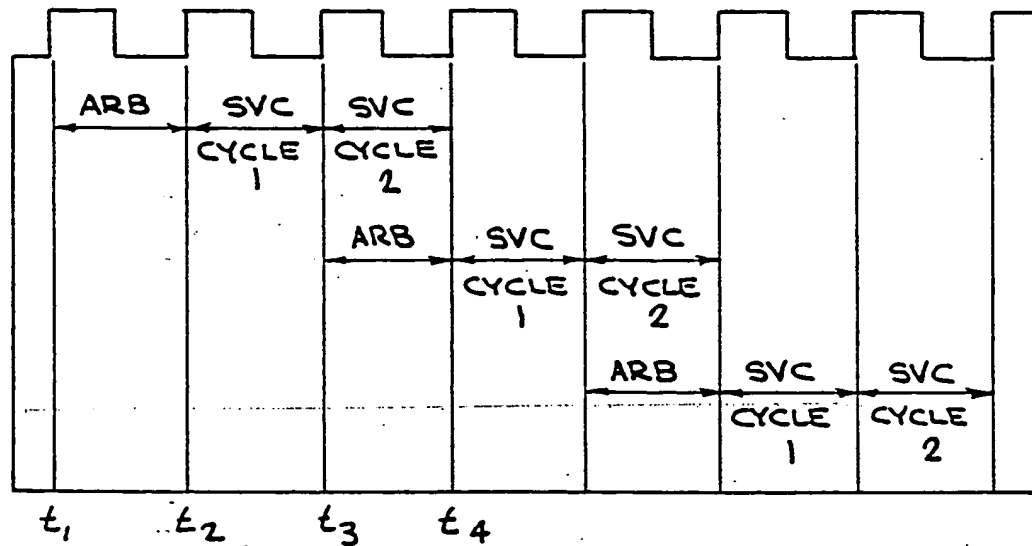


FIG. 15

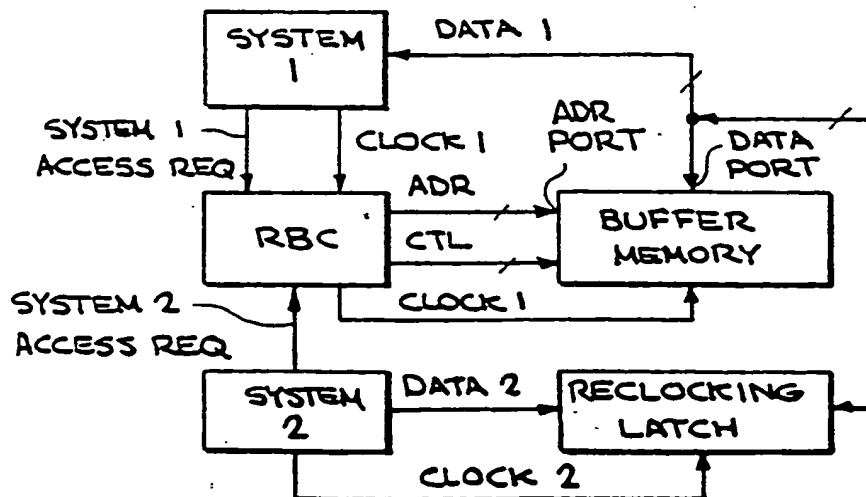


FIG. 16

12

EUROPEAN PATENT APPLICATION

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22 Date of filing: 03.03.87

30 Priority: 06.03.86 US 836936

43 Date of publication of application:
14.10.87 Bulletin 87/42

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AT BE CH DE ES FR GB GR IT LI LU NL SE

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30.05.90 Bulletin 90/22

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54 Addressing arrangement for a RAM buffer controller.

57 There is disclosed herein a RAM buffer controller for managing the address input lines of a RAM buffer to simulate the operation of two FIFO's therein. There is also disclosed apparatus for allowing random access by a node processor in a local area network node using the RAM buffer controller to manage transmit and receive FIFO's to have random access to any address in the address space of the buffer without restriction to FIFO boundaries. There is also disclosed apparatus for transmitting packets from said buffer organized into one or two linked lists. Further, there is disclosed apparatus for allowing independent initialization of any of the pointers in the RAM buffer controller which are not currently selected, and for allowing software requests for read or write access by the node processor. Further, there is disclosed apparatus and a method for recording status and length information at the end of a packet instead of in front thereof and for allowing any incoming packet to be flushed without saving status information or to be flushed while saving its status information.

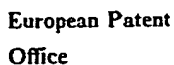
EP 0 241 129 A3
 A3



DOCUMENTS CONSIDERED TO BE RELEVANT			
Category	Citation of document with indication, where appropriate, of relevant passages	Relevant to claim	CLASSIFICATION OF THE APPLICATION (Int. Cl. 4)
X	WO-A-8 400 836 (WESTERN ELECTRIC) * Abstract; page 9, line 9 - page 10, line 3; page 14, lines 1-24; figure 4 *	1,2,8	G 06 F 5/06
Y		3-7,10, 34,38, 41	
A		17,29, 30,43- 47,39, 17-20	
X	US-A-4 334 287 (WIEDENMAN et al.) * Claim 1; figure 2 *	1,8,14, 15,23, 26,27, 31	
Y		3-7,37, 38,41, 42	
A		9,11-13 ,16,21, 22,24, 25,28, 30,32, 40	
			TECHNICAL FIELDS SEARCHED (Int. Cl.4)
			G 06 F
The present search report has been drawn up for all claims			
Place of search THE HAGUE		Date of completion of the search 01-03-1990	Examiner COHEN B.
CATEGORY OF CITED DOCUMENTS X : particularly relevant if taken alone Y : particularly relevant if combined with another document of the same category A : technological background O : non-written disclosure P : intermediate document T : theory or principle underlying the invention E : earlier patent document, but published on, or after the filing date D : document cited in the application L : document cited for other reasons & : member of the same patent family, corresponding document			



DOCUMENTS CONSIDERED TO BE RELEVANT			
Category	Citation of document with indication, where appropriate, of relevant passages	Relevant to claim	CLASSIFICATION OF THE APPLICATION (Int. Cl. 4)
Y	IEEE, PROC. OF THE ANNUAL SYMPOSIUM ON COMPUTER ARCHITECTURE, 23rd-25th March 1977, Long Beach, CA, US, pages 159-164; PARKER et al.: "Hardware/software tradeoffs in a variable word width, variable queue length buffer memory" * Abstract; page 159, column 2, lines 1-45; page 161, column 2, lines 26-14; figures 1,2 *	5,6,7	
Y	FR-A-2 507 372 (NIPPON TELEGRAPH) * Claim 1 *	10	
X	WO-A-8 002 755 (BURROUGHS) * Page 1, line 18 - page 2, line 26; figure 1; page 6, lines 16-28 *	33	
Y		34,35,36	
A		8	
Y	DE-A-2 629 498 (SIEMENS) * Claim 1 *	35-38,41,42	TECHNICAL FIELDS SEARCHED (Int. Cl.4)
Y	US-A-3 493 935 (QUESTA et al.) * Abstract; column 4, line 65 - column 5, line 13; figure 2 *	42	
A		17-20,43-47	
	---	-/-	
The present search report has been drawn up for all claims			
Place of search THE HAGUE		Date of completion of the search 01-03-1990	Examiner COHEN B.
CATEGORY OF CITED DOCUMENTS X : particularly relevant if taken alone Y : particularly relevant if combined with another document of the same category A : technological background O : non-written disclosure P : intermediate document T : theory or principle underlying the invention E : earlier patent document, but published on, or after the filing date D : document cited in the application L : document cited for other reasons & : member of the same patent family, corresponding document			



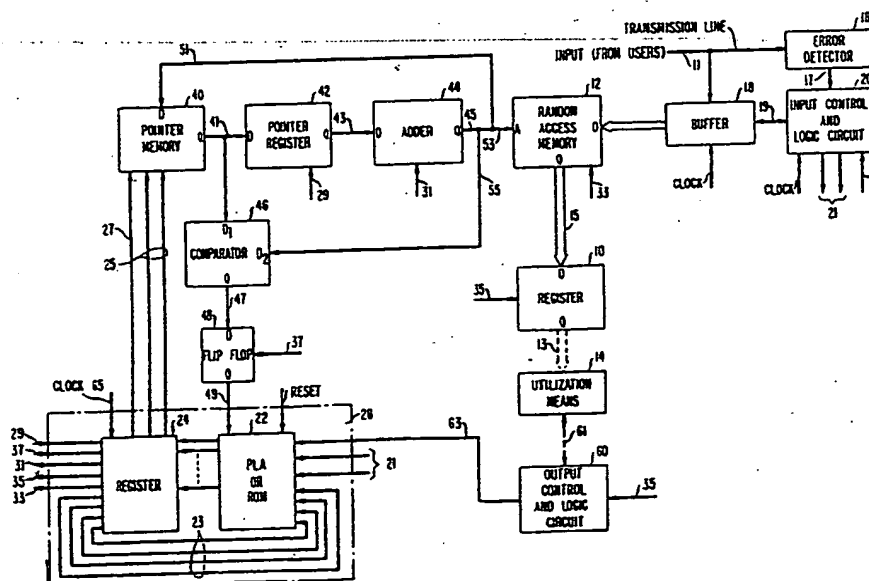
EP 87 30 1850

DOCUMENTS CONSIDERED TO BE RELEVANT			CLASSIFICATION OF THE APPLICATION (Int. Cl. 4)
Category	Citation of document with indication, where appropriate, of relevant passages	Relevant to claim	
A	IBM TECHNICAL DISCLOSURE BULLETIN, vol. 24, no. 1A, June 1981, pages 365-367, Armonk, NY, US; GOLDSTEIN: "Enqueue-top for recovery" * Whole document *	10	TECHNICAL FIELDS SEARCHED (Int. Cl. 4)
A	WO-A-8 400 835 (WESTERN ELECTRIC) * Page 2, lines 23-37; page 3, lines 26-31 *	17-20, 43-47	
A	IEEE TRANSACTIONS ON NUCLEAR SCIENCE, vol. NS-32, no. 1, February 1985, pages 105-107; APPEL et al.: "A high capacity FASTBUS multiple event buffer" * Abstract; page 106, column 1, line 30 - page 107, column 1, line 28; figures 2,3 *	17-20, 43-47	
The present search report has been drawn up for all claims			
Place of search THE HAGUE		Date of completion of the search 01-03-1990	Examiner COHEN B.
CATEGORY OF CITED DOCUMENTS X: particularly relevant if taken alone Y: particularly relevant if combined with another document of the same category A: technological background O: non-written disclosure P: intermediate document		T: theory or principle underlying the invention E: earlier patent document, but published on, or after the filing date D: document cited in the application L: document cited for other reasons &: member of the same patent family, corresponding document	

INTERNATIONAL APPLICATION PUBLISHED UNDER THE PATENT COOPERATION TREATY (PCT)

<p>(51) International Patent Classification³ : G06F 13/00</p>	<p>A1</p>	<p>(11) International Publication Number: WO 84/ 0083: (43) International Publication Date: 1 March 1984 (01.03.84)</p>
<p>(21) International Application Number: PCT/US83/00190 (22) International Filing Date: 15 February 1983 (15.02.83) (31) Priority Application Number: 407,877 (32) Priority Date: 13 August 1982 (13.08.82) (33) Priority Country: US (71) Applicant: WESTERN ELECTRIC COMPANY, INC. [US/US]; 222 Broadway, New York, NY 10038 (US). (72) Inventor: FRASER, Alexandre, Gibson ; 62 Carriage House Road, Bernardsville, NJ 07924 (US). (74) Agents: HIRSCH, A., E., Jr. et al.; Post Office Box 901, Princeton, NJ 08540 (US).</p>		<p>(81) Designated States: DE (European patent), FR (European patent), GB (European patent), JP, NL (European patent), SE (European patent). Published <i>With international search report.</i> <i>With amended claims.</i></p>

(54) Title: FIRST-IN, FIRST-OUT (FIFO) MEMORY CONFIGURATION FOR QUEUE STORAGE



(57) Abstract

A first-in, first-out queue has a random access memory (RAM) for storing a plurality of information words, serialism. A controller is used to insure that only after a complete message, comprising the information words, has been received will a word of that message be read out. Three pointers are used to effect this result. A read pointer addresses the location in the RAM from where a word may be read. A write pointer addresses the location in the RAM where a word may be entered. A third pointer addresses the location in the RAM where the last word of a complete message is stored.

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FIRST-IN, FIRST-OUT (FIFO) MEMORY
CONFIGURATION FOR QUEUE STORAGE

Technical Field

This invention relates to digital communications systems and, in particular, to FIFO memories enabling multiword messages to be written therein one word at a time and to be read out therefrom as complete messages.

Background of the Invention

In telecommunications transmission systems, there arises sometimes a need to store information from a transmitter because the receiver is busy. The information, however, must be recovered from storage in the same order in which transmitted. This process is known as first-in, first-out (FIFO).

In U. S. Patent No. 3,979,733 granted September 7, 1976 to Mr. A. G. Fraser, there was disclosed a FIFO queue. A read register, in the aforesaid Fraser patent, has recorded therein the address of the next memory cell to be read. Likewise, a write register has recorded therein the address of the next memory cell available for storing data. A comparison of the read and write registers indicates whether the memory cells are all full, all empty or partially empty.

It is frequently desired to use a FIFO queue as a communication path between a data producing process and a separate data consuming process. These processes are typically independent of one another and may not even be controlled from a common clock source. That is, the two processes may be asynchronous with respect to one another. The path carries messages from the data producing process to the data consuming process and it is frequently required to guarantee that the consuming process obtain only complete messages.

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A problem arises if the data-producing process has to abandon a message after inserting some of it into the FIFO queue. In that case, the consuming process may already have started the message. This situation arises, for example, when the data producing process is a transmission line with error detection equipment and the consuming process is a computer. If an error is detected partially through an arriving message, it is usually desired that the message be discarded and not processed by the computer.

In the aforesaid Fraser patent, however, there is no way of knowing when a complete message has been received. In the absence of such knowledge it is common practice to use two memories. The first memory is used to assemble one complete message at a time. The second memory is used to operate as a FIFO system in which the unit of storage is a message. Such an arrangement usually requires high speed processing circuitry which then becomes a bottleneck limiting throughput. Further, because two memories are used, circuitry is necessarily duplicated.

Summary of the Invention

In accordance with the illustrative embodiment of the present invention there is disclosed a first-in, first-out memory system comprising a controller for storing, seriatim, a plurality of information words in a memory device. Only when a complete message has been stored will the controller allow the aforesaid information words to be read-out in the first-in, first-out sequence.

More particularly, the storage device is a random access memory (RAM) with a capacity for storing N information words, where N is an integer. Such a memory is known as a cyclic memory. The controller comprises three pointer registers: W, R, and L. The W and R pointer registers hold addresses of the RAM locations previously accessed. The L pointer register holds the address of the RAM location where a complete message ends.

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When an information word is to be entered in the RAM for storage, the W pointer register is incremented to $W+1$ modulo N and the address $W+1$ is compared with the address in the R pointer register. If the values of the two addresses are not equal, the information word is entered in the storage location in the RAM whose address is $W+1$. Thereafter, the address in the W pointer register is changed inside the controller to $W+1$.

When an information word is to be read from the RAM, the address in the R pointer register is compared with the address in the L pointer register. If the values of the two addresses are not equal, the address in the R pointer register is incremented by one to $R+1$ modulo N . The information word, addressed by the R pointer, namely, the location in the RAM addressed by $R+1$, is read from the RAM. Thereafter, the address in the R pointer register inside the controller is updated to indicate the new address $R+1$.

When a complete message has been entered in the RAM, the address of the location where the message ends is entered in the L pointer register. That is, the address in the W pointer register is copied into the L pointer register, thereby effectively moving the L pointer to address the location in the RAM where the complete message ends.

When an error in transmission is detected during the entry of information words in the memory and part way through a message, the previously entered partial message must be discarded. This result is achieved by copying the address in the L pointer register into the W pointer register.

An advantage of the present invention is the use of a single random access memory for the first-in, first-out queue and a single control circuit to implement the necessary queue management. Furthermore, very high speed components are not required and there no longer arises a processing bottleneck. Additionally, because only

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complete messages are read-out, synchronization problems are avoided.

Brief Description of the Drawings

FIG. 1 shows a block diagram of a first-in, first-out queue for storing information words and for reading out the words only when a complete message has been stored;

FIG. 2 shows a block diagram of a message received by the first-in, first-out queue; and

FIG. 3 shows a diagram illustrating the use of the memory in FIG. 1.

Detailed Description

Referring to FIG. 1, there is shown a random access memory (RAM) 12 for storing information received over a transmission line 11. The information may be transmitted as a message comprising a plurality of information words. One such message is shown in FIG. 2. The format of a message may vary depending on the particular use. The message in FIG. 2 comprises a header 32, data 34, and end of message flag 36. The information words are stored in RAM 12, in FIG. 1, in the order in which they are received over line 11. After all the information words in a complete message have been entered in RAM 12, they are available for retrieval by a utilization means 14, such as a digital computer, a switching machine, and the like. When the message is retrieved from RAM 12, the first information word to be entered therein is read out first. That is, a first-in, first-out queue is realized using a random access memory, RAM 12.

As messages of data arrive, they are entered simultaneously into error detector 16 and buffer 18. If there exists no error resulting from transmission, lead 17 is enabled. Simultaneously, enabling input control and logic circuit 20 receives information about the beginning and end of messages from buffer 18 to be described hereinbelow.

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The use of RAM 12 may be explained by referring to FIG. 3. RAM 12 may be thought of as a circular storage device where information words, comprising a variable number of bits may be stored. Thus, if the number of information words that may be stored in RAM 12 is denoted by N, the information words are entered sequentially in locations 0,1,2 ... N-1. After the location with address N-1 has been filled, the next location to be filled will have address 0. That is, successive locations can be addressed using modulo N arithmetic.

A pointer R addresses the location of the last word that was read from RAM 12. When a word is to be read out from RAM 12, pointer R is incremented by 1 modulo N and the contents of that word at location R+1 is read.

Likewise, a pointer W addresses the location where a word of information was last entered. Thus, when a word of information is to be entered into RAM 12, pointer W is incremented by 1 and the word is entered in the location whose address is W+1 modulo N.

The locations between R+1 modulo N and W represent the words of information available to be read from RAM 12. According to the present invention, messages comprising a plurality of words of information are entered in RAM 12 when received but will not be available for being read from RAM 12 until all information in a message is entered therein. This is achieved by the use of a third pointer L.

When a complete message is received, the end of message flag 36, FIG. 2, carrying a special code is interpreted in the input control circuit 20, FIG. 1. This condition is transmitted over bus 21 to a programmable logic array (PLA) or a read only memory (ROM) 22. Referring to FIG. 3 again, when a complete message is received, the contents of pointer W is copied into pointer L. That is, pointer L addresses the location of the last information word in the complete message.

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Before a word is read from RAM 12, pointers R and L are compared. If they point to the same location, the queue does not have a complete message, and a word cannot be read from RAM 12. If pointers R and L address
5 different locations, pointer R is incremented by 1 and the contents of location R+1 modulo N is read from RAM 12.

Because the memory is circular and is a random access memory, it is necessary to prevent destruction of valid data by insuring that pointer W will not be advanced
10 beyond pointer R. According to the present invention, pointer W is maintained so that when incremented by 1, W+1 modulo N will never be equal to pointer R. That is, there is provided a cushion of one word location.

Referring again to FIG. 1, there is shown a
15 pointer memory 40 which holds three addresses: pointers W, L, and R. In the illustrative embodiment, each pointer is twelve bits long. Pointer register 42 is large enough to hold a single pointer. Adder 44 is designed to add either zero or one modulo N to its input value, namely, an
20 address indicated by pointer W, L or R. RAM 12 has one word location for every information word from buffer 18. Further RAM 12 has as many locations as can be addressed by a single pointer. Thus, in the illustrative
embodiment, because each pointer has twelve bits, each
25 pointer can address 4,096 words (2^{12}) in RAM 12, that is, $N = 4,096$, in this illustration. Register 10 is wide enough to store a single word of information read from RAM 12. Comparator 46 compares the values of two
addresses and produces an output. If the two addresses
30 are equal, the output from comparator 46 is one, otherwise, the output therefrom is zero. Flip-flop 48 stores the output from comparator 46 and that output is made available over lead 49 to ROM 22.

The pointers W, R and L, as stated hereinbefore,
35 address locations in RAM 12 so that information may be written therein or read therefrom. Pointers W and R move

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cyclicly through RAM 12, one location at a time. As stated hereinbefore, RAM 12 may be conceptualized as a circular buffer having three pointers, W, R and L.

Pointer memory 40 is addressed by leads 25. The address, comprising two bits, indicates the pointer W, R, or L. Depending on the pointer location addressed, a pointer appears at the Q port and on bus 41. When a pointer, which appears on bus 51, is to be entered in pointer memory 40, the location is indicated by address leads 25 and the control lead 27 is enabled. The pointer on bus 41 appears at the D port of pointer register 42 and at the D₁ port of comparator 46.

The pointer on bus 41 at the D port is copied into pointer register 42 by enabling the control lead 29. The contents of pointer register 42 is always present at its Q port and on bus 43.

As stated hereinabove, adder 44 adds either zero or one modulo N to its input. The pointer, from pointer register 42, appears at the D port of adder 44 and the quantity, either a zero or a one, appearing on the input lead 31, are added. The resulting sum, pointer +0 (or 1) modulo N, appears at the output port Q and on bus 45. Bus 45 branches into three separate buses: 51, to the D port of pointer memory 40; 53, to the address port A of the RAM 12; and 55, to the D₂ port of comparator 46.

Words which are to be entered in RAM 12 appear at the D port thereof. The pointer indicating the address of the memory cell into which the word is to be entered appears at the A port thereof. When the control lead 33 is enabled, the word at the D port is entered into the memory cell of RAM 12 addressed by the pointer on bus 53.

Words which are to be read from RAM 12 are addressed by the pointer on bus 53 again through the A port. Thereafter, the words appear at the Q port of

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RAM 12 and via bus 15 at the D port of register 10. When control lead 35 is enabled, the word which was read from RAM 12 is entered into register 10 and appears at the Q port thereof and on bus 13.

5 The input values at the D_1 and D_2 ports of comparator 46 are compared. If the input values are equal, the output is one; otherwise, the output is zero. The output from comparator 46 is transmitted via bus 47 to the D port of flip-flop 48. When control lead 37 is
10 enabled, the output from comparator 46 is entered in flip-flop 48. The value of flip-flop 48 is continuously present on lead 49 to the ROM 22.

ROM 22 and control register 24 together form a control circuit 26. When utilization means 14 is ready to
15 receive a message, a signal is transmitted over bus 61 to output control and logic circuit 60. Thereafter, output control circuit 60 issues a read command over lead 63 to the ROM 22. Depending on the status of the eight leads 21, 23, 49 and 63 at the input to ROM 22, an
20 instruction is read therefrom and is transferred to control register 24 simultaneously with a clock pulse on lead 65. The contents of register 24, namely the instructions from ROM 22, define the state of the device during the next clock period. Thus, the register 24 holds
25 the current values for the twelve leads 23, 25, 27, 29, 31, 33, 35 and 37. The leads 23 carry a number, comprising four bits, which is fed back as an input to PLA 22 in order that the following state may be generated. The next state as defined by the next instruction from
30 ROM 22 depends upon both the previous state identified by the number on leads 23 and the new inputs on control leads 21, 49 and 63.

Reading Information from the FIFO Queue

In response to a command from utilization
35 means 14, during a first clock period, control circuit 26 causes the address of the R pointer to be transmitted over leads 25 to the pointer memory 40. Simultaneously

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therewith, control lead 29 is enabled thereby causing the aforesaid R pointer to be entered in pointer register 42. In a second clock period, the lead 31 carries the value zero so that the R pointer is passed intact through

5 adder 44 and appears at the D_2 port of comparator 46. During the same second clock period, leads 25 carry the address of the L pointer to pointer register 40 and the L pointer appears at the D_1 port of comparator 46. The output from comparator 46 is entered in flip-flop 48 by
10 enabling lead 37. If the queue is empty or does not have a complete message, pointers L and R are equal and the value on lead 49 is one. If the queue, however, has either a partially read message or at least one complete message, the value on lead 49 is zero.

15 During the third clock period, the value on lead 31 is a one and the adder 44 is allowed to settle down. The value of the output, $R+1$, from adder 44 appears on bus 53 and points to the word to be read from RAM 12.

20 During the fourth clock period, if the value on lead 49 was zero during the second clock period, the lead 35 is enabled and a word, which was read from RAM 12 during the third clock period, is entered into register 10. The word is continuously available at the Q port of register 10 and on bus 13. Lead 35 also carries
25 an input signal to the output control logic circuit 60 where it signifies when one word has been read.

During the aforesaid third clock period the incremented R pointer, namely $R+1$, appears on bus 51 at the D port of pointer memory 40. During the fourth clock
30 period, leads 25 carry the address of the R pointer and lead 27 is enabled, thereby causing the incremented R pointer, $R+1$, on bus 51 to be entered as the new R pointer in pointer memory 40.

Writing Information into the FIFO Queue

35 In response to a word being received on line 11, input control circuit 20 transmits this status over leads 21 to control circuit 26. Thereafter, during the

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first clock period, leads 25 carry the address of the W pointer in pointer memory 40 and the W pointer thus addressed is entered in pointer register 42 by enabling lead 29. During the second clock period, the lead 31 carries the value one so that the output from adder 44 is W+1 which appears at the D₂ port of comparator 46.

Simultaneously, the leads 25 carry the address of the R pointer in pointer memory 40 and the R pointer thus addressed appears at the D₁ port of comparator 46.

During the third clock period, the value of the signal on the lead 31 remains at one and the output from comparator 46 is entered in flip-flop 48 by enabling lead 37. If the output from comparator 46, namely, the contents of flip-flop 48, is zero, it means the R pointer and W+1 are not equal and the word at the D port of RAM 12 may be entered therein.

Thus, during the fourth clock period, if the R pointer and W+1 are not equal, the word at the D port of RAM 12 is entered in the location addressed by W+1 at the A port thereof by enabling lead 33. Lead 33 also carries an input signal to the input control circuit 20 and signifies when a word from buffer 18 has been entered in RAM 12.

During the same fourth clock period, leads 25 carry the address of the W pointer in pointer memory 40 and lead 27 is enabled so that the incremented value, W+1, of the W pointer on bus 51 may be entered in the pointer memory 40.

Signalling the Entry of a Complete Message

When a complete message has been received, as indicated by the end of message flag 36, of FIG. 2, leads 21, in FIG. 1, carry this status to the control circuit 26. In response thereto, the L pointer is updated. This updating requires two clock periods.

During the first clock period, leads 25 carry the address of the W pointer in pointer memory 40. The lead 29 is enabled and the W pointer is copied into the

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pointer register 42. During the second clock period lead 31 carries the value zero. Thus, the W pointer appears on bus 51. The leads 25 carry the address of the L pointer in pointer memory 40. Lead 27 is enabled
5 thereby copying the value of the W pointer on bus 51 into the L pointer.

Resetting the FIFO Queue

The FIFO queue is reset, or initialized, to render the queue empty by copying the value of the
10 R pointer into the L and W pointer locations in pointer memory 40. This process requires three clock periods.

During the first clock period, the leads 25 carry the address of the R pointer in pointer memory 40 and the R pointer is entered in pointer register 42 by
15 enabling lead 29.

During the second clock period, the value on lead 31 is zero so that the R pointer appears on bus 51. The leads 25 carry the address of the L pointer and by enabling lead 27 the value of the R pointer is copied into
20 the L pointer location in pointer memory 40.

Likewise, during the third clock period, the value on lead 31 remains at zero so that the R pointer continues to appear on bus 51. The leads 25 carry the address of the W pointer and by enabling lead 27 the value
25 of the R pointer is copied into the W pointer location.

Contention Resolution

If both write command leads 21 and read command lead 63 are simultaneously enabled, the control circuit 26 gives priority to the write command leads 21, even if the
30 FIFO queue is full and the operation cannot be completed. The next operation should be the read command on lead 63.

Input Control Process

As stated hereinabove, when an information word is received it is placed in buffer 18 and in error
35 detector 16. With the complete entry of each word in buffer 18, input control circuit 20 transmits a command over leads 21 to ROM 22 thereby causing the information

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word in buffer 18 to be transferred therefrom and entered into RAM 12.

As stated hereinabove, after each information word is entered in error detector 16, the presence or
5 absence of a transmission error is detected therein. If no transmission error had been detected, lead 17 is enabled. Meanwhile, each word entered in buffer 18 is compared with an end of message flag. When an end of message is detected, that condition is transmitted over
10 bus 19 to input control circuit 20.

When lead 17 is enabled and an end of message signal is received over bus 19 from buffer 18, the input control circuit 20 transmits a command over leads 21 to ROM 22. In response thereto, the address in the W pointer
15 is copied into the L pointer, thereby signifying the entry of a complete message in RAM 12. Thereafter, this message may be read from RAM 12.

Should an error in transmission be detected in error detector 16, however, lead 17 will not be enabled.
20 In the absence of a signal on lead 17, a different command is transmitted from input control circuit 20 over leads 21 to ROM 22. In response thereto, the address in the L pointer is copied into the W pointer, thereby returning the W pointer to an initial position. As a result, no
25 erroneous information words are retained in RAM 12.

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Claims

1. A first-in, first-out memory system comprising
- 5 means for storing a message comprising a plurality of information words, and means for controlling
- 1) the entry of said information words in said storage means, and
- 2) the reading of said information words from
- 10 said storage means in the same order in which said information words were entered therein, only after all said information words of said message have been stored.
2. The first-in, first-out memory system according to claim 1 wherein said memory is a random access
- 15 memory.
3. The first-in, first-out memory system according to claim 1 wherein said means for controlling the entry of said information words comprises a first pointer for addressing the location in said storage means where
- 20 said information word is to be stored.
4. The first-in, first-out memory system according to claim 1 wherein said means for controlling the reading of said information words comprises a second pointer for addressing the location in said storage means
- 25 from where said information word is to be read.
5. The first-in, first-out memory system according to claim 4 wherein said means for controlling the entry of said information words and the reading of said information words further comprises a third pointer for
- 30 addressing the location in said storage means where the last of said information words in said complete message is stored.
6. The first-in, first-out memory system according to claim 5 further comprising
- 35 means for discarding said message or a part thereof before one of said information words is read.



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7. A first-in, first-out queue comprising
a random access memory for storing information
words, and

a controller for permitting the read-out of
5 said words after a complete message comprising a plurality
of said words has been stored in said random access memory,
and said controller comprises a programmable logic array
(PLA) or read only memory (ROM) and a register.

8. The first-in, first-out queue according to
10 claim 7 wherein said controller further comprises a pointer
memory for storing three pointers.

9. The first-in, first-out queue according to
claim 9 wherein said controller further comprises
a pointer register for storing any one of said
15 pointer,

an adder for either adding to said pointer
from said pointer register the number one, thereby
incrementing the value of said pointer from said pointer
register, or adding to said pointer from said pointer
20 register the number zero,

a comparator having two input signals thereto,
a first input signal being the value of said pointer read
from said pointer memory and a second input signal being an
output signal from said adder, and
25 means for storing the resulting output signal
from said comparator.

10. A first-in, first-out memory system for
multiword messages comprising

an addressable cyclic memory
30 a W pointer register for writing words into
said memory at the address of the W pointer,
an R pointer register for reading words from
said memory at the address of the R pointer,
an L pointer register for identifying the last
35 word of a multiword message,

means for incrementing the W pointer register,
and comparing the W pointer register to the R pointer

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register before storing a new message word and only if the incremented W pointer and the R pointer are not the same,

means for comparing the contents of the R pointer register with the contents of the L pointer

5 register, and incrementing the R pointer register and reading the memory word addressed by the R pointer register only if the R and L pointers are not the same,

means responsive to an end of message signal for copying the contents of the W pointer register into the

10 L pointer register, and

means responsive to an error in a new memory word for copying the contents of the L pointer register into the W pointer register.



AMENDED CLAIMS

(received by the International Bureau on 31 August 1983 (31.08.83))

1. (Original claims 1 thru 5 amended) A first-in, first-out memory system comprising means for storing a message comprising a plurality of information words, and
5 means for controlling the entry of the information words in the storage means, and
the reading of the information words from the storage means in the same order in which the information words were entered, only after all the information words of the message have been stored,
the means for controlling the entry of the information words comprises a first pointer for addressing
10 the location in the storage means where the information word is to be stored, and
the means for controlling the reading of the information words comprises a second pointer for addressing the location in the storage means from where the
15 information word is to be read,

CHARACTERIZED IN THAT

~~the means for controlling the entry of the~~
information words and the reading of the information words further comprises a third pointer for addressing the
25 location in the storage means where the last of the information words in the complete message is stored.

2. (Original claim 6 amended) First-in, first-out memory system according to claim 1,

CHARACTERIZED BY

30 means for discarding the message or a part thereof before one of the information words is read.

3. (Original claims 7 & 8 amended) First-in, first-out queue comprising

35 a random access memory for storing information words, and

a controller for permitting the read-out of the words after a complete message comprising a plurality



of the words has been stored in the random access memory, and the controller comprises a programmable logic array (PLA) or read only memory (ROM) and a register.

CHARACTERIZED IN THAT

5 the controller further comprises a pointer memory for storing three pointers.

4. (Original claim 10 amended) First-in, first-out memory system for multiword messages comprising an addressable memory

10 a W pointer register for writing words into the memory at the address of the W pointer,

an R pointer register for reading words from the memory at the address of the R pointer, means for incrementing the W pointer register, and comparing the W
15 pointer register to the R pointer register,

CHARACTERIZED BY

the addressable memory being cyclic, means for storing a new message word and only if the incremented W pointer and the R pointer are not the same,

20 an L pointer register for identifying the last word of a multiword message,

means for comparing the contents of the R pointer register with the contents of the L pointer register and incrementing the R pointer register and
25 reading the memory word addressed by the R pointer register only if the R and L pointers are not the same,

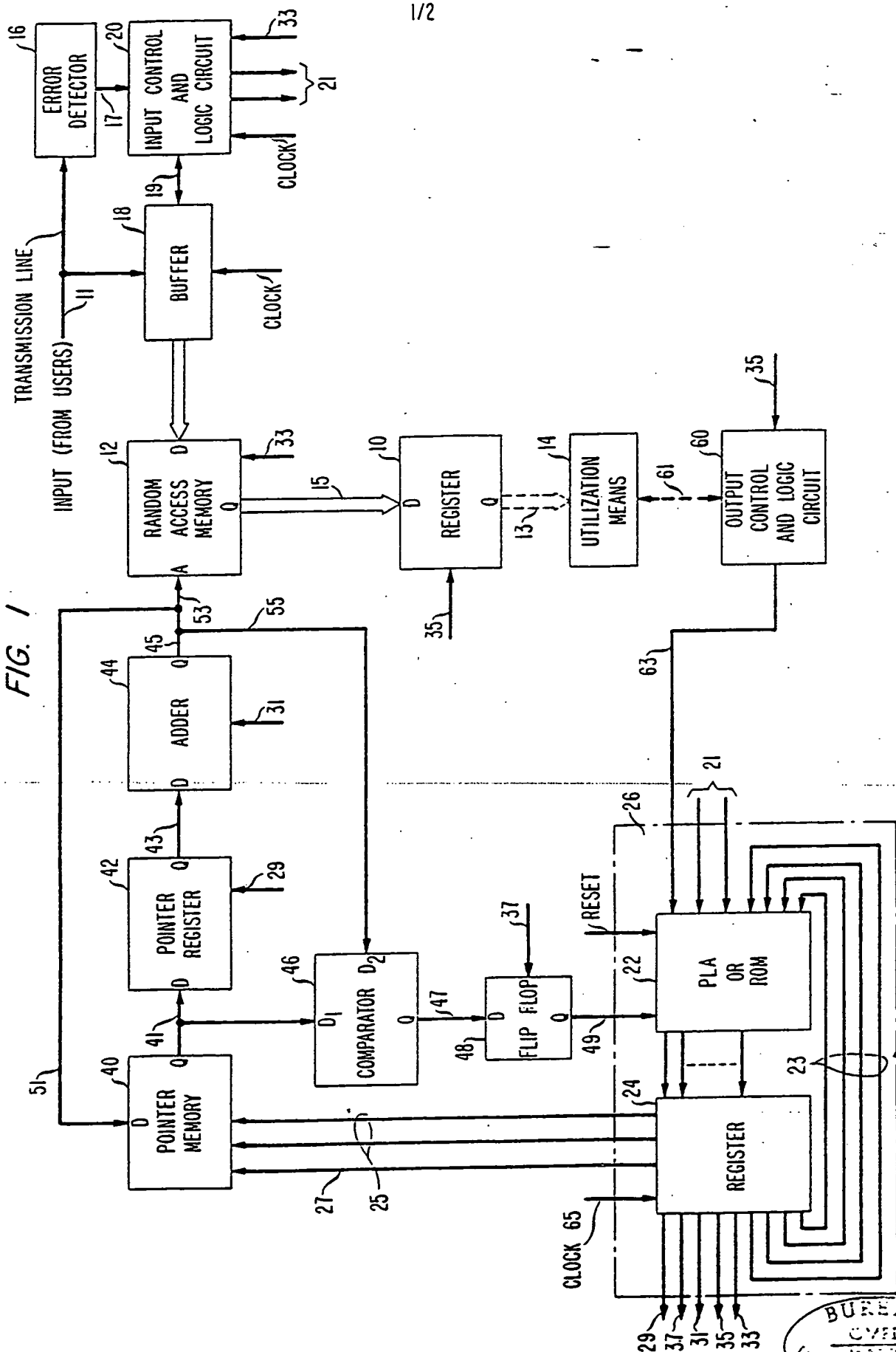
means responsive to an end of message signal for copying the contents of the W pointer register into the L pointer register, and

30 means responsive to an error in a new memory word for copying the contents of the L pointer register into the W pointer register.



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FIG. 1



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FIG. 2

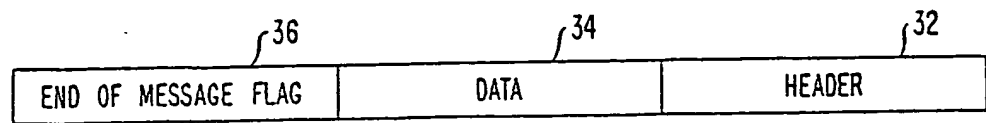
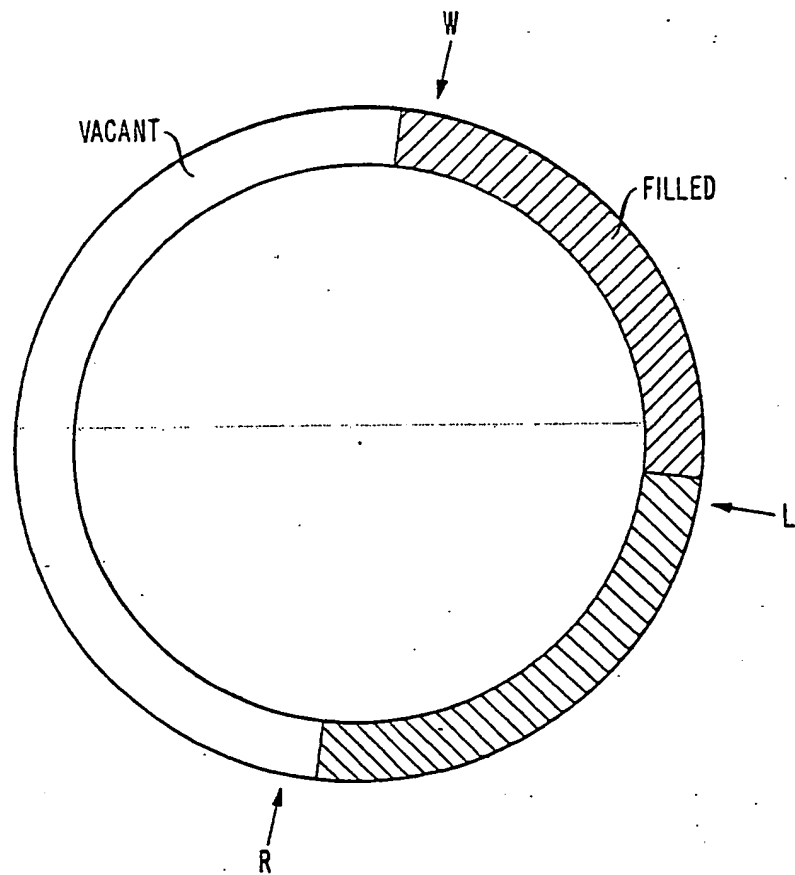


FIG. 3



INTERNATIONAL SEARCH REPORT

International Application No PCT/US 83/00190

I. CLASSIFICATION OF SUBJECT MATTER (If several classification symbols apply, indicate all) ³ According to International Patent Classification (IPC) or to both National Classification and IPC IPC ³ : G 06 F 13/00						
II. FIELDS SEARCHED <div style="text-align: center; border-top: 1px solid black; border-bottom: 1px solid black;">Minimum Documentation Searched ⁴</div> <table style="width: 100%; border-collapse: collapse;"> <tr> <td style="width: 30%; border-bottom: 1px solid black;">Classification System</td> <td style="border-bottom: 1px solid black;">Classification Symbols</td> </tr> <tr> <td style="padding: 5px;">IPC³</td> <td style="padding: 5px;">G 06 F 13/00; H 04 L 11/20</td> </tr> </table> <div style="text-align: center; border-top: 1px solid black; border-bottom: 1px solid black;">Documentation Searched other than Minimum Documentation to the extent that such Documents are Included in the Fields Searched ⁵</div>			Classification System	Classification Symbols	IPC ³	G 06 F 13/00; H 04 L 11/20
Classification System	Classification Symbols					
IPC ³	G 06 F 13/00; H 04 L 11/20					
III. DOCUMENTS CONSIDERED TO BE RELEVANT ¹⁴						
Category ⁶	Citation of Document, ¹⁴ with indication, where appropriate, of the relevant passages ¹⁷	Relevant to Claim No. ¹⁸				
X	IBM Technical Disclosure Bulletin, vol. 24, no. 3, published August 1981 (New York, US) Cary et al.: "Data buffer management", pages 1502-1503 <div style="text-align: center;">--</div>	1-7,9				
Y	FR, A, 2260141 (HONEYWELL BULL) 29 August 1979 see figure 3; page 5, line 13 - page 13 <div style="text-align: center;">--</div>	1-5,7,9				
Y	US, A, 3818461 (WARD et al.) 18 June 1974 see figures 2,3; column 8, line 6 - column 9, line 69 <div style="text-align: center;">--</div>	1-5,7,9				
A	IBM Technical Disclosure Bulletin, vol. 20, no. 8, published January 1978 (New York, US) Chapman et al.: "Data management in a circular buffer", pages 3309-3310 <div style="text-align: center;">--</div>	1-4,7-10 ./.				
<div style="display: flex; justify-content: space-between;"> <div style="width: 45%;"> <p>⁹ Special categories of cited documents: ¹⁶</p> <p>"A" document defining the general state of the art which is not considered to be of particular relevance</p> <p>"E" earlier document but published on or after the international filing date</p> <p>"L" document which may throw doubts on priority claim(s) or which is cited to establish the publication date of another citation or other special reason (as specified)</p> <p>"O" document referring to an oral disclosure, use, exhibition or other means</p> <p>"P" document published prior to the international filing date but later than the priority date claimed</p> </div> <div style="width: 45%;"> <p>"T" later document published after the international filing date or priority date and not in conflict with the application but cited to understand the principle or theory underlying the invention</p> <p>"X" document of particular relevance; the claimed invention cannot be considered novel or cannot be considered to involve an inventive step</p> <p>"Y" document of particular relevance; the claimed invention cannot be considered to involve an inventive step when the document is combined with one or more other such documents, such combination being obvious to a person skilled in the art.</p> <p>"A" document member of the same patent family</p> </div> </div>						
IV. CERTIFICATION						
Date of the Actual Completion of the International Search ³ <div style="text-align: center; font-size: 1.2em;">15th June 1983</div>		Date of Mailing of this International Search Report ³ <div style="text-align: center; font-size: 1.2em;">15 JUL 1983</div>				
International Searching Authority ¹ <div style="text-align: center; font-weight: bold;">EUROPEAN PATENT OFFICE</div>		Signature of Authorized Officer ²⁰ <div style="text-align: center;"> G.L.M. Gruydenberg </div>				

III. DOCUMENTS CONSIDERED TO BE RELEVANT (CONTINUED FROM THE SECOND SHEET)		
Category ¹⁶	Citation of Document, ¹⁶ with indication, where appropriate, of the relevant passages ¹⁷	Relevant to Claim No ¹⁸
A	US, A, 3553651 (BIRD et al.) 5 January 1971 ---	1-4,7,9
A	US, A, 2907004 (CHIEN et al.) 29 September 1959 ---	1-4,7,9
A	US, A, 3633173 (EDGE) 4 January 1972 ---	1
A	US, A, 3979733 (FRASER) 7 September 1976 (cited in the application) -----	1

ANNEX TO THE INTERNATIONAL SEARCH REPORT ON

INTERNATIONAL APPLICATION NO.

PCT/US 83/00190 (SA 4859)

This Annex lists the patent family members relating to the patent documents cited in the above-mentioned international search report. The members are as contained in the European Patent Office EDP file on 11/07/83

The European Patent Office is in no way liable for these particulars which are merely given for the purpose of information.

Patent document cited in search report	Publication date	Patent family member(s)	Publication date
FR-A- 2260141	29/08/75	None	
US-A- 3818461	18/06/74	US-A- 3748652	24/07/73
		DE-A, B, C 2317687	25/10/73
		GB-A- 1418986	24/12/75
		CA-A- 978656	25/11/75
		JP-A- 49011035	31/01/74
US-A- 3553691	05/01/71	None	
US-A- 2907004		None	
US-A- 3633173	04/01/72	DE-A- 2111153	07/10/71
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		NL-A- 7604729	11/11/76
		FR-A, B 2310595	03/12/76
		DE-A, B, C 2620220	18/11/76
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		AU-A- 1366176	10/11/77
		CA-A- 1056065	05/06/79
		JP-A- 51138103	29/11/76
		SE-C- 420256	14/01/82
		SE-A- 7605192	10/11/76

For more details about this annex :
see Official Journal of the European Patent Office, No. 12/82